

NEW LOWER BOUNDS ON UNIPROCESSOR SCHEDULABILITY

		Implicit deadlines ($d = p$)	Constrained deadlines ($d \leq p$)	Arbitrary deadlines (d, p unrelated)
FP	Arbitrary utilization			
	Utilization bounded by a constant c			
EDF (Feasibility)	Arbitrary utilization			
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EDF (Feasibility)	Arbitrary utilization	Polynomial- time algorithm	Exponential- time algorithm	Exponential- time algorithm
	Utilization bounded by a constant c	Polynomial- time algorithm	Pseudo-poly- time algorithm for $0 < c < 1$	Pseudo-poly- time algorithm for $0 < c < 1$

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	Utilization bounded by a constant c	Polynomial-time algorithm for $c \leq \ln 2$ and RM^\dagger priorities	Pseudo-poly.-time algorithm for $0 < c < 1$	Exponential-time algorithm for $0 < c < 1$
EDF (Feasibility)	Arbitrary utilization	Polynomial-time algorithm	Weakly coNP-complete (E&R, SODA'10)	Weakly coNP-complete (E&R, SODA'10)
	Utilization bounded by a constant c	Polynomial-time algorithm	Pseudo-poly.-time algorithm for $0 < c < 1$	Pseudo-poly.-time algorithm for $0 < c < 1$

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EDF (Feasibility)	Arbitrary utilization	Polynomial-time algorithm	Strongly coNP-complete (ECRTS'15)	Strongly coNP-complete (ECRTS'15)
	Utilization bounded by a constant c	Polynomial-time algorithm	Pseudo-poly.-time algorithm for $0 < c < 1$	Pseudo-poly.-time algorithm for $0 < c < 1$

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EDF (Feasibility)	Arbitrary utilization	Polynomial-time algorithm	Strongly coNP-complete (ECRTS'15)	Strongly coNP-complete (ECRTS'15)
	Utilization bounded by a constant c	Polynomial-time algorithm	Weakly coNP-complete for $0 < c < 1$ (RTSS'15)	Weakly coNP-complete for $0 < c < 1$ (RTSS'15)

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FP	Arbitrary utilization	Weakly NP-complete (New)	Weakly NP-complete (New)	Exponential-time algorithm
	Utilization bounded by a constant c	Polynomial-time algorithm for $c \leq \ln 2$ and RM [†] priorities	Pseudo-poly-time algorithm for $0 < c < 1$	Exponential-time algorithm for $0 < c < 1$
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FP	Arbitrary utilization	Weakly NP-complete (New)	Weakly NP-complete (New)	Weakly NP-hard
	Utilization bounded by a constant c	Polynomial-time algorithm for $c \leq \ln 2$ and RM [†] priorities	Weakly NP-complete for $0 < c < 1$ (New)	Weakly NP-hard for $0 < c < 1$
EDF (Feasibility)	Arbitrary utilization	Polynomial-time algorithm	Strongly coNP-complete (ECRTS'15)	Strongly coNP-complete (ECRTS'15)
	Utilization bounded by a constant c	Polynomial-time algorithm	Weakly coNP-complete for $0 < c < 1$ (RTSS'15)	Weakly coNP-complete for $0 < c < 1$ (RTSS'15)

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