Ch 4: Trees

"it's a jungle out there..."

"I think that I will never see a linked list useful as a tree; Linked lists are used by everybody, but it takes real smarts to do a tree"



Trees: examples (*corporate structure*)



Trees: examples (your Unix file system)



Definitions

- A tree t is a <u>finite nonempty</u> set of elements.
 One of these elements is called a root, and the remaining elements (if any) are partitioned into trees that are called <u>subtrees</u> of t.
- children of the root -- root of subtree
- parent of an element/ grandparent/ ancestor/ descendent/...
- leaves: elements with no children
- degree of an element: # children
- degree of a tree: max degree of any element

















Binary trees

- A binary tree t is a finite (possibly empty) collection of elements. When the binary tree is not empty, it has a root element and the remaining elements (if any) are partitioned into two binary trees, called the left subtree and the right subtree of t
- Notes:
 - a binary tree may be empty
 - each element has exactly two (perhaps empty) subtrees
 - the subtrees/ children are ordered

Example binary tree: an *expression tree*





Properties of binary trees

- A binary tree of n elements has ?? edges
 - exactly n-1
- The height/ depth of a binary tree is the number of levels in it.
 A binary tree of height h has at least ?? and at most ??? elements
 - at least h
 - at most 2^h 1
- The height of a binary tree with n elements is at most ?? and at least ???
 - at most <mark>n</mark>
 - at least [log₂ (n+1)]

Binary trees (review from yesterday)

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Properties of binary trees (review from yesterday)

- A binary tree of n elements has n-1 edges
- The height/ depth of a binary tree is the number of levels in it.
 A binary tree of height h has at least h and at most 2^h - 1 elements
- The height of a binary tree with n elements is at most n and at least [log₂ (n+1)]

Binary trees: some more definitions



A FULL binary tree (Ht h ==> #-elements = 2^{h} -1)

Binary trees: some more definitions



A COMPLETE binary tree (levels fill in from the left) Representing binary trees As an array:

Use the canonical ordering -- node with canonical ordering i goes into T[i]

T[i]'s left child at ?? T[i]'s right child at ?? T[i]'s parent is at ??

Problems with this approach: memory inefficient (works fine for full or complete binary trees)

Representing binary trees As a pointer to a treeNode:



Representing binary trees As a pointer to a treeNode:

template <class T>
class node{
public:

};

T data; node <T> * LC; node <T> * RC;

LC	data	RC
----	------	----

Representing binary trees As a pointer to a treeNode:

```
templ ate <cl ass T>
cl ass node{
publ i c:
    node(T x, node<T>* t1=NULL, node<T>* t2=NULL);
    T data;
    node <T> * LC;
    node <T> * RC;
};
```

```
{data = x; LC = t1; RC = t2;}
```









Representing binary trees

template <class C>
class node{
public:
 node(C x, node<C>* t1, node<C> * t2);
 C data;
 node<C> * LC;
 node<C> * RC;
 };



Recursive programming on binary trees: counting elements

```
// in the user (client) program
int countElements(node<char> * T)
{
    if (T == NULL) return 0;
    return (
        1 // itself
        + countElements(T->LC)
        + countElements(T->RC) );
}
```

Recursive programming on binary trees: computing depth

```
// in the user (client) program
int depth(node<char> * T)
{
    if (T == 0) return 0;
    return (
        1 +
        max (depth(T->LC),
            depth(T->RC)) );
}
```

Recursive programming on binary trees: comparing trees

Recursive programming on binary trees: comparing trees

```
// in the user (client) program
bool what(node<int> * T1, node<int> * T2)
{
    if (T1 == NULL) return (T2 == NULL);
    if (T2 == NULL) return false;
    return ( what(T1->LC, T2->LC)
        &&
        what(T1->RC, T2->RC) );
}
```

Recursive programming on binary trees: preorder traversal

```
// in the user (client) program
void preOrder(node<char> * T)
{
    if (T == 0) return;
        cout << T->data;
        preOrder(T->LC);
        preOrder(T->RC);
}
```

Recursive programming on binary trees: inorder traversal

```
// in the user (client) program
void inOrder(node<char> * T)
{
    if (T == 0) return;
        inOrder(T->LC);
        cout << T->data;
        inOrder(T->RC);
}
```

Recursive programming on binary trees: postorder traversal

```
// in the user (client) program
void postOrder(node<char> * T)
{
    if (T == 0) return;
        postOrder(T->LC);
        postOrder(T->RC);
        cout << T->data;
}
```

Recursive programming on binary trees: ???

```
// in the user (client) program
treeNode<char * mystery(node<char> * T)
{
    if (T == 0) return NULL;
    return new node<char>(
        T->data,
        mystery(T->LC),
        mystery(T->RC)
        );
}
    mystery returns a copy
```

Recursive programming on binary trees: ???

```
// in the user (client) program
treeNode<char * puzzle(node<char> * T)
{
    if (T == 0) return NULL;
    return new node<char>(
        T->data,
        puzzle(T->RC),
        puzzle(T->LC)
        );
}
puzzle returns a mirror-image
```