

Generating Realistic TCP Workloads

Felix Hernandez-Campos Ph. D. Candidate Dept. of Computer Science Univ. of North Carolina at Chapel Hill

Recipient of the 2001 CMG Fellowship

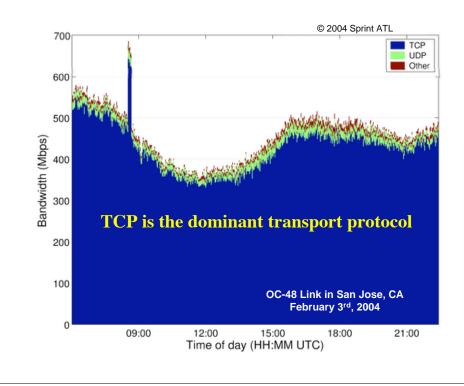
Joint work with F. Donelson Smith and Kevin Jeffay

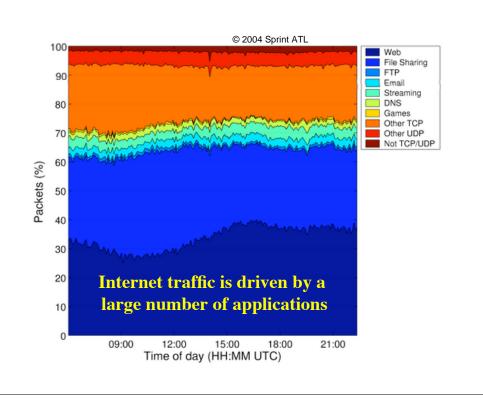


Experimental Networking Research and Performance Evaluation

- Evaluating network protocols and mechanisms requires careful experimentation
 - -Network simulation (NS, Opnet, etc.)
 - -Network testbeds
- A critical element of these experiments is the **traffic workload**

- Are current workloads *realistic*?
- Let's look at some measurements - Sprint's tier-1 backbone network

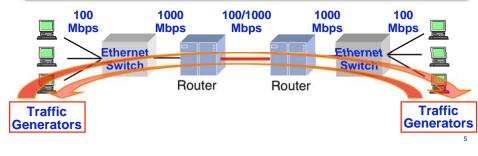






Internet Traffic Generation **Testbed example**

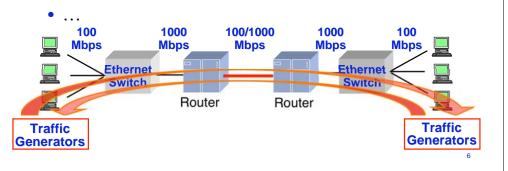






Internet Traffic Generation Testbed example

- Evaluate queuing mechanisms in the routers
- Evaluate transport protocols
- Evaluate intrusion/anomaly detection mechanisms
- Evaluate traffic monitoring techniques





Traffic Generation

State-of-the-Art

Open-loop

- -Large number of sophisticated models
 - » Packet-level modeling
- -But TCP is a closed-loop protocol
 - » Open-loop traffic generation breaks reliability, flow control, and congestion control

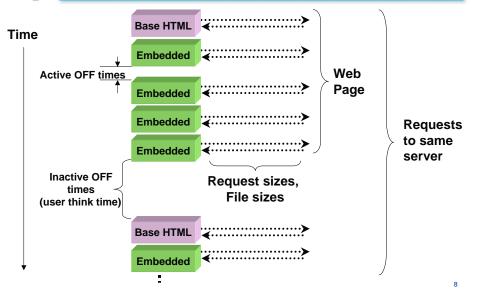
Closed-loop

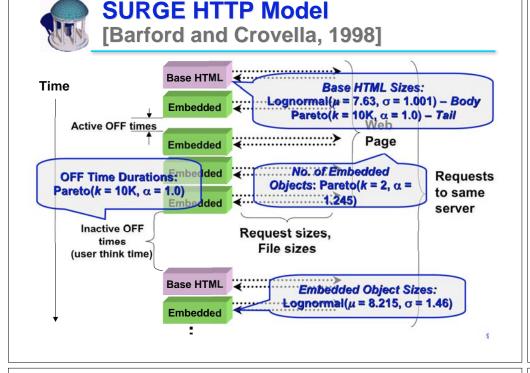
- The idea is to simulate the behavior of users/applications
 - » Source-level modeling of specific application



7

Application-Specific Modeling HTTP Model Example







Application-Specific Models Shortcomings

- Creating application models is a challenging, timeconsuming task
 - Traffic mixes are driven by a large number of applications
- Set of dominant applications evolves quickly
 - Applications themselves also evolve
 - We cannot even identify a significant fraction of the traffic
- Modeling closed application protocols requires reverse-engineering
- Privacy considerations complicate data acquisition
 - TCP header traces are ok, application data is not



Abstract source-level modeling

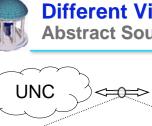
- Application-neutral technique to describe the source-level behavior of any TCP connection
- Efficient analysis applicable any arbitrary trace of TCP headers (no analysis of payloads)
- We can also measure network-level parameters, such as round-trip times, receiver window sizes, etc.

Source-level trace replay

- Replaying abstract source-level behavior

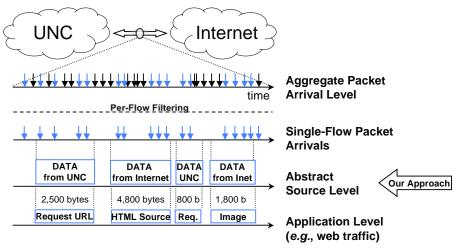
Validation in network testbed

- We wrote a distributed, scalable traffic generator (tmix)
- Comparison of original and synthetic traces



Different Views of Internet Traffic

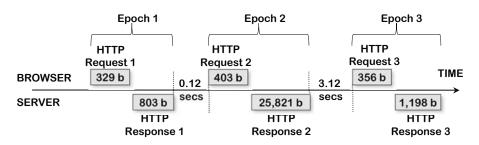
Abstract Source-level Modeling





Client-Server Applications



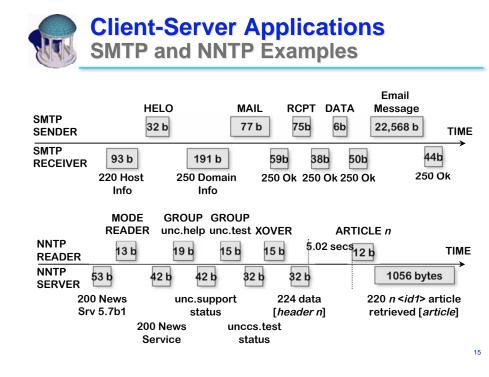


- We call a pair of application data units (ADUs) that carry a request/response exchange an *epoch*
- Quiet times are also part of the workload of TCP



• Abstract source-level model for describing the workload of TCP connections

- Each connection is summarized using a *connection vector* of the form $C_i = (e_1, e_2, ..., e_n)$ with $n \ge 1$ epochs
 - Each epoch has the form $e_j = (a_j, ta_j, b_j, tb_j)$
- Connection vectors can be extracted from trace of TCP segment headers
 - Sequence number directionality, timing analysis, write size and packet size interactions
 - $-O(n \log n) + O(n^*W)$

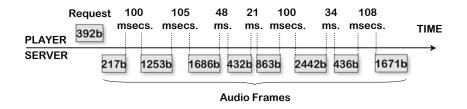




13

Beyond the Client-Server Model

Icecast – Internet Radio



- Server PUSH applications do not follow the traditional client-server model
- The sequential a-b-t model is still applicable – Make a_i and tb_i zero

16

Beyond the Client-Server Model NNTP in Stream-Mode

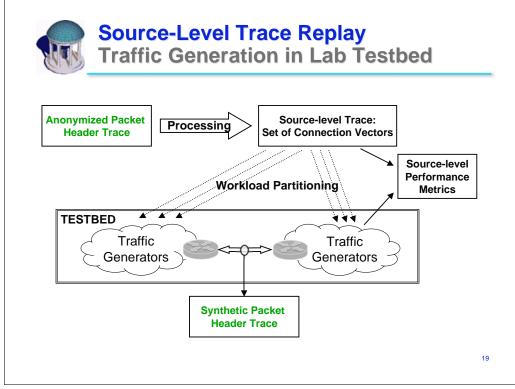
| MODE CHEC STREAM <id12 PEER 13 b 41 b</id12 | > <id3></id3> | TAKETHIS < <i>id2</i> > [<i>article</i>] 15,678 bytes | |
|---|---|---|------|
| | 438 238 4 | 1 b 49 b TI 38 438 m't don't | ME |
| | send < <i>id2</i> > se < <i>id1</i> > < <i>i</i> d | end send d <i>3</i> > < <i>id4</i> > | |
| Protocol Bitfield Interest | red Piece Piece <i>i j</i> 16397 b 16397 b | Request Piece Piece Piece M 100 16397 b | ١. |
| PEER B 68 b 657 b 5b BitTorrent Interested Protocol | Piece j | 17b 16397 b Request Piece k m | TIME |
| Bitfield | Request Piece i | Request Piece / | 17 |

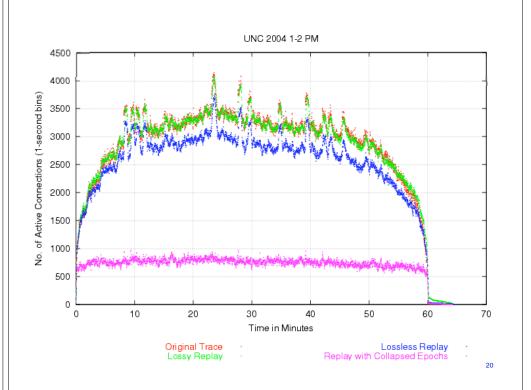


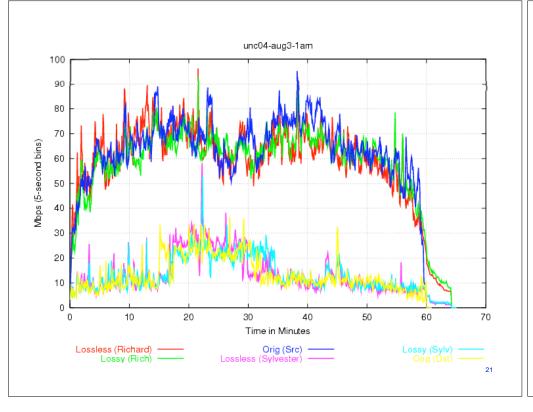
- Some connections are said to exhibit *data exchange concurrency*
- Two reasons:
 - Increasing performance
 - Enabling natural concurrency
- Concurrent a-b-t model describes each side of the connection separately

 $((a_1, ta_1), (a_2, ta_2), \dots, (a_n, ta_n))$ $((b_1, tb_1), (b_2, tb_2), \dots, (b_m, tb_m))$

Concurrency can be detected with high probability
p.seqno > *q.ackno* and *q.seqno* > *p.ackno* - O(n*W)









• New method for modeling traffic mixes

- Empirically-derived connection vectors
- Studied sequential vs. concurrent dichotomy
- Fully automated, efficient analysis
- New traffic generation approach
 - Enables comparison of real and synthetic traffic
 - Implemented a distributed traffic generator
 - Techniques for scaling traffic load

