The UNIVERSITY *of* NORTH CAROLINA
at CHAPEL HILL

Rate-Based Resource Allocation Models for Multimedia Computing and Embedded Systems

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http://www.cs.unc.edu/Research/dirt

Rate-Based Resource Allocation
The case against static priority scheduling

- Static priority scheduling in general, and Rate Monotonic scheduling in particular, dominates in the real-time systems literature
  - VxWorks, VRTX, QNX, pSOSystems, LynxOS all support static priority scheduling
- Does one size fit all?
  - "When you have a hammer, everything looks like a nail"
- Problems with static priority scheduling
  - Feasibility is dependent on a predictable environment and well-behaved tasks.

Rate-Based Resource Allocation
Overview

- The problem:
  - How to allocate resources in an environment wherein…
    - » Work arrives at well-defined but highly variable rates
    - » Tasks may exceed their execution time estimates
  - … and still guarantee adherence to deadlines
- The thesis:
  - Static priority scheduling is the wrong tool for the job (existing task models are too simplistic)
  - Rate-based scheduling abstractions can simplify the design and implementation of many real-time systems and improve performance and resource utilization

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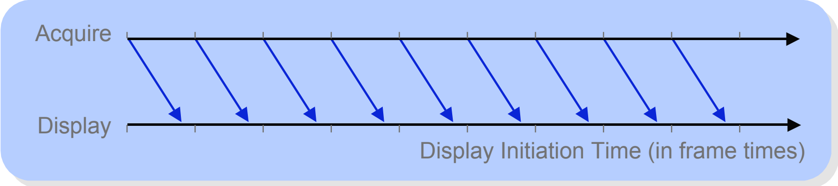
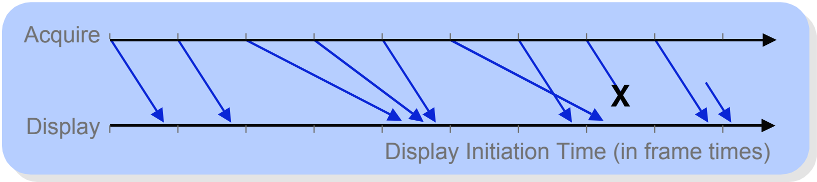
The Case Against Priority Scheduling
Example: Display-side multimedia processing



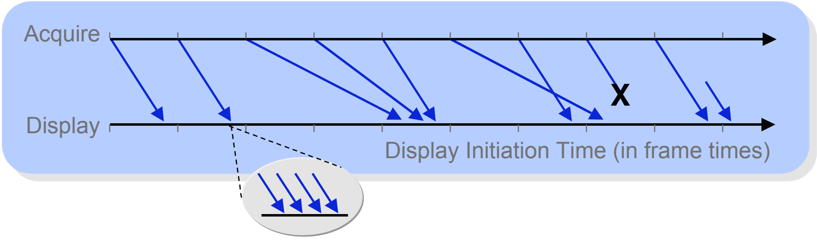
- The problem: Receive frames from the network and deliver to a display application so as to ensure...
  - Continuous playout
  - Minimal playout latency
- The theory: Multimedia is easy — it's periodic!
  - Apply existing theory of periodic or sporadic tasks

Display-side Media Processing
The practice



- Nothing is periodic in a distributed system!
- The effects of distributed systems pathology:
  - Variable message transmission times
  - Out-of-order message arrivals
  - Lost & duplicate messages

Display-side Media Processing
Managing the Network Interface



- Packets fragmented in the network must be reassembled
  - *Messages* have deadlines, *packets* do not
  - Applications know about messages, operating systems do not

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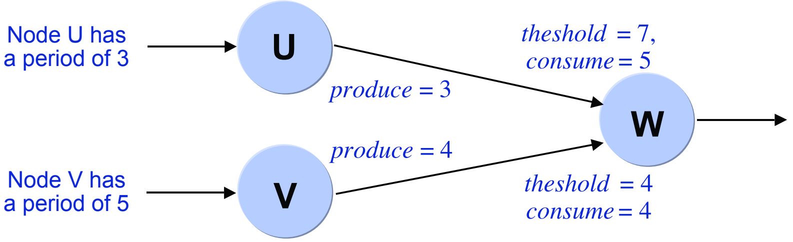
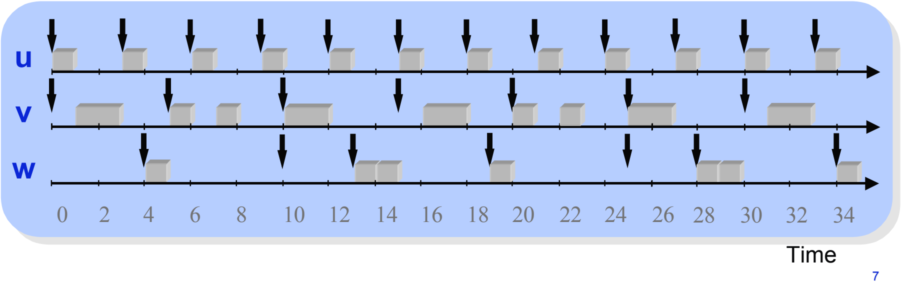
The Case Against Priority Scheduling
Example: Signal processing data flow graphs





Rate-Based Computing
Approaches

- Extend the Liu and Layland model of real-time tasks to allow the expression of real-time rates
  - Hierarchical "server-based" scheduling — Create a "server" process that is scheduled as a periodic task and internally schedules the processing of aperiodic events
  - Event-based scheduling — Process aperiodic events as if they were generated by a virtual "well behaved" periodic process
- Adapt "fluid-flow" models of resource allocation developed in the networking community for bandwidth allocation to CPU scheduling
  - Provide a "virtual processor" abstraction wherein each task logically executes on a dedicated processor with $1/f(n)$ the capacity of the physical processor

An Event-Based Rate Model

The Rate-Based Execution (*RBE*) model

- Tasks make progress at the rate of processing x events every y time units and each event is processed within d time units (in the best case)

- For task i with rate specification (x_i, y_i, d_i), the j^{th} event for task i, arriving at time $t_{i,j}$, will be processed by time

$$D(i,j) = \begin{cases} t_{i,j} + d_i & \text{if } 1 \le j \le x_i \\ \mathrm{MAX}(t_{i,j} + d_i,\ D(i, j-x_i)+y_i) & \text{if } j > x_i \end{cases}$$

  - $D(i,j)$ gives the earliest possible deadline for the j^{th} instance of task i ($\ge t_{i,j} + d_i$)

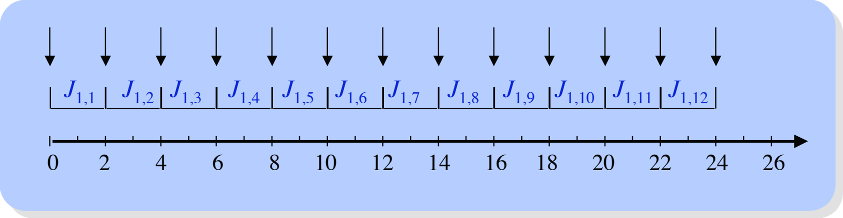
The RBE Task Model

Example: Periodic arrivals, periodic service

- Task with rate specification ($x = 1$, $y = 2$, $d = 2$)

$$D(i,j) = \begin{cases} t_{i,j} + d_i & \text{if } 1 \le j \le x_i \\ \mathrm{MAX}(t_{i,j} + d_i,\ D(i, j-x_i)+y_i) & \text{if } j > x_i \end{cases}$$

  - Deadlines separated by at least $y = d = 2$ time units
  - Deadlines occur at least 2 time units after a job is released



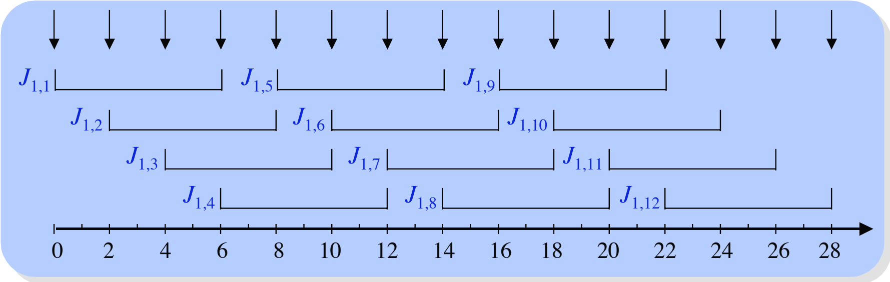
The RBE Task Model

Example: Periodic arrivals, *deadline ≠ period*

- Task with rate specification ($x = 1$, $y = 2$, $d = 6$)

$$D(i,j) = \begin{cases} t_{i,j} + d_i & \text{if } 1 \le j \le x_i \\ \mathrm{MAX}(t_{i,j} + d_i,\ D(i, j-x_i)+y_i) & \text{if } j > x_i \end{cases}$$

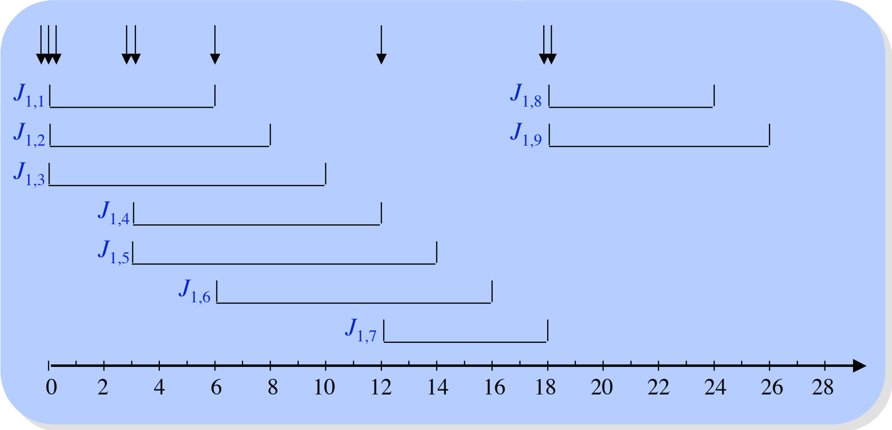
  - Deadlines separated by at least $y = 2$ time units and occur at least $d = 6$ time units after a job is released



The RBE Task Model

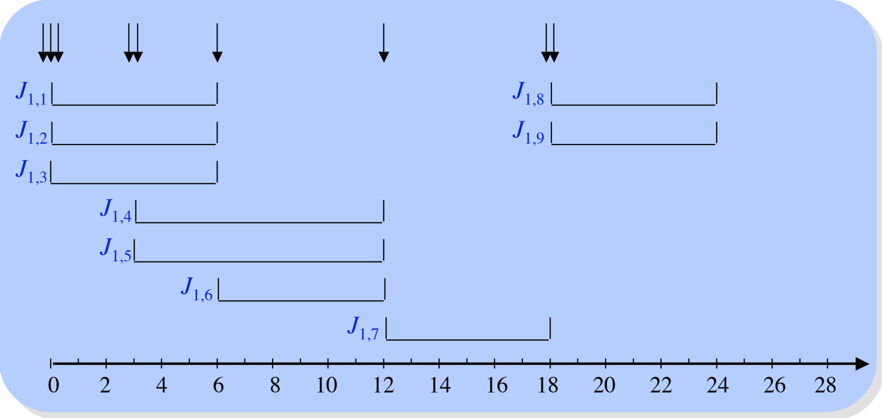
Bursty arrivals

- Task with rate specification ($x = 1$, $y = 2$, $d = 6$)
  - Deadlines separated by at least $y = 2$ time units and occur at least $d = 6$ time units after a job is released



The RBE Task Model
Bursty arrivals

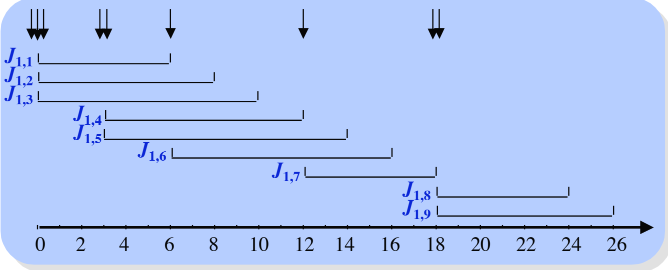
- Task with rate specification ($x = 3$, $y = 6$, $d = 6$)
  - Deadlines separated by at least $y = 6$ time units and occur at least $d = 6$ time units after a job is released



The RBE Task Model
Comparison of rate specifications

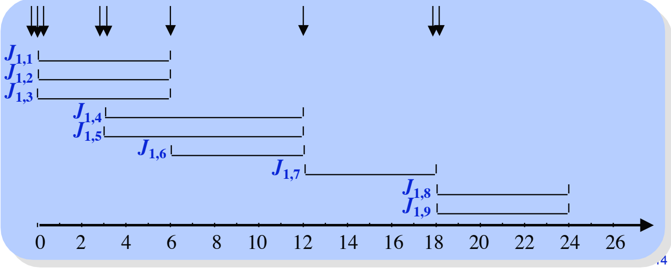
Rate specification

(x=1, y=2, d=6)



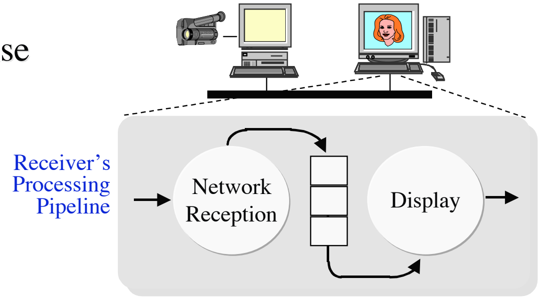
Rate specification

(x=3, y=6, d=6)

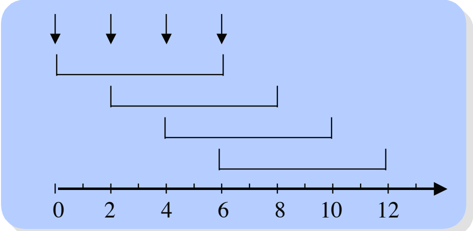


The RBE Task Model
RBE features/properties

- Provides better response time for non-real-time activities by integrating application-level buffering with the system run queue

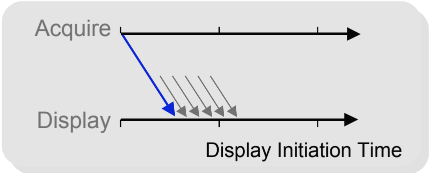


Rate specification
(x = 1, y = 2, d = 6)

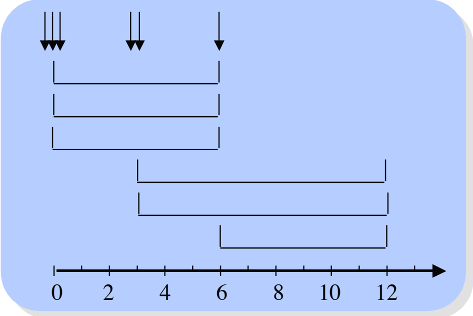


The RBE Task Model
RBE features/properties

- Provides a more natural way of modeling inbound packet processing of fragmented messages

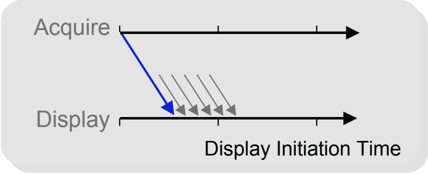


Rate specification
(x = 3, y = 6, d = 6)

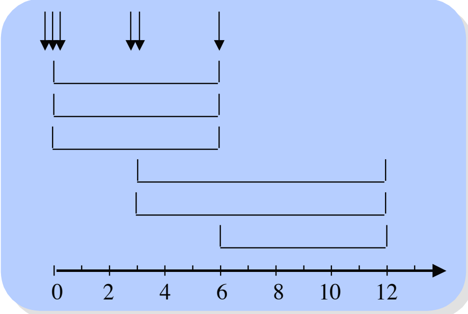


The RBE Task Model
RBE features/properties

- Provides isolation from arrival rates that exceed the rate specification
  - (But does not provide isolation from tasks exceeding their stated execution time)



Rate specification
($x = 3$, $y = 6$, $d = 6$)



Fluid Flow Resource Allocation
Proportional share resource allocation

- Tasks are allocated a *share* of the processor's capacity
  - Task *i* is assigned a *weight* w_i
  - Task *i*'s *share* of the CPU at time *t* is

$$f_i(t) = \frac{w_i}{\sum_{j \in A(t)} w_j}$$

- If tasks' weights remain constant in $[t_1, t_2]$ then task *i* receives

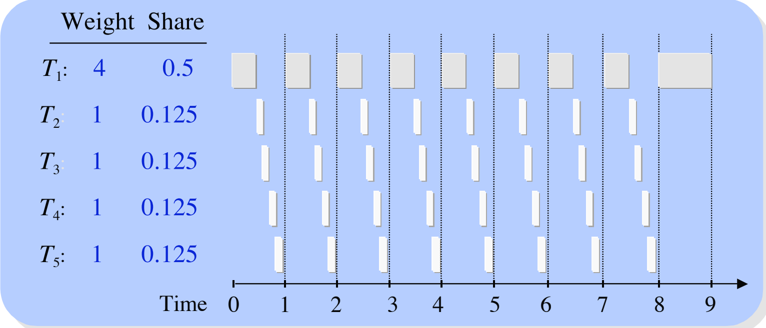
$$S_i(t_1,t_2) = \int_{t_1}^{t_2} f_i(t)\, dt = \frac{w_i}{\sum_j w_j}(t_2 - t_1)$$

units of execution time in $[t_1, t_2]$

Proportional Share Resource Allocation
Fluid scheduling example

- Weighted round robin scheduling with an infinitesimally small quantum
- In $[t_1, t_2]$ (if total weight doesn't change) T_i receives

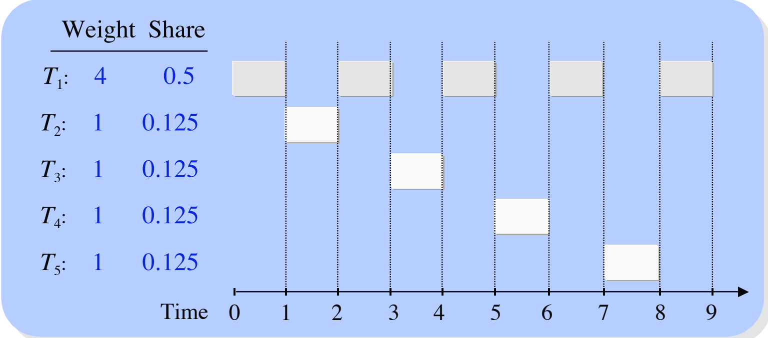
$$S_i(t_1,t_2) = \int_{t_1}^{t_2} f_i(t)\, dt = \frac{w_i}{\sum_{j \in A(t)} w_j}(t_2 - t_1)$$



	Weight	Share
T_1:	4	0.5
T_2:	1	0.125
T_3:	1	0.125
T_4:	1	0.125
T_5:	1	0.125

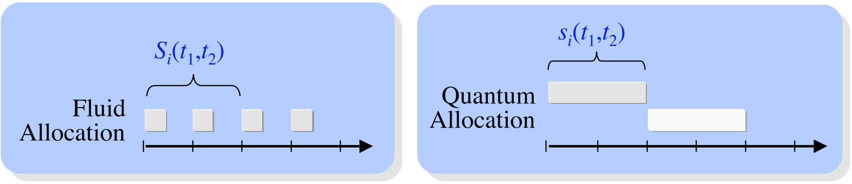
Proportional Share Resource Allocation
Quantum scheduling example

- Weighted round robin scheduling with integer quanta
  - $q = 1$
- The quantum system doesn't proportionally allocate the resource over all time intervals



	Weight	Share
T_1:	4	0.5
T_2:	1	0.125
T_3:	1	0.125
T_4:	1	0.125
T_5:	1	0.125

Proportional Share Resource Allocation
Task scheduling metrics & goals



- Schedule tasks so that their performance is as close as possible to that in the *fluid* system
- Why is fluid allocation important?
  - What about real-time allocation?!

Approximating Fluid Allocation
Why is this so important?

- Fluid allocation implies real-time progress
- Weights are used to allocate a *relative* fraction of the CPU's capacity to a task

$$f_i(t) = \frac{w_i}{\Sigma_j w_j}$$

- Real-time progress requires a *constant* fraction of the CPU's capacity

$$\forall t,\ f_i(t) = \textit{execution cost}_i \times \textit{execution frequency}_i$$

  - If a task must execute for 16 *ms* every 33 *ms* then allocating $f = 0.5$ ensures real-time execution
- Thus real-time performance can be achieved by adjusting weights dynamically so that the share remains constant

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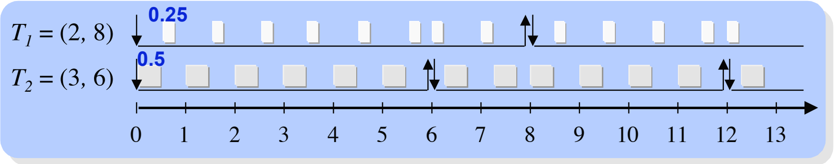
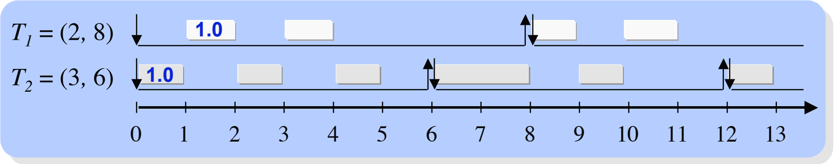
Proportional Share Resource Allocation
Real-time scheduling example

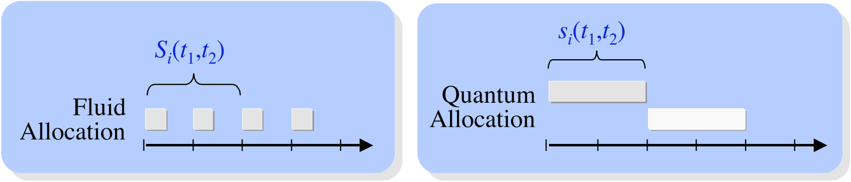
- Periodic tasks allocated a share equal to their processor utilization
  - Round-robin scheduling with infinitesimally small quantum



  - With unit-sized quantum



Proportional Share Resource Allocation
Task scheduling metrics & goals



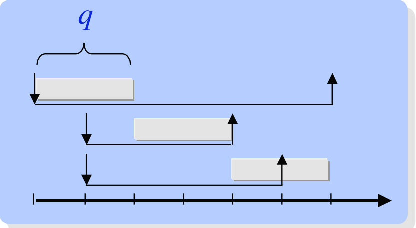
- Goal: Schedule tasks so that their performance is as close as possible to that in the *fluid* system
- Define the allocation error for task *i* at time *t* as

$$lag_i(t) = \left(\begin{array}{c}\textit{allocation the task would have}\\ \textit{received in the fluid system}\end{array}\right) - \left(\begin{array}{c}\textit{allocation the task has received}\\ \textit{in the quantum system}\end{array}\right)$$

$$= S_i(t_i,t) - s_i(t_i,t)$$

- Schedule tasks so that the lag is bounded for all tasks over all time intervals
  - What is the least upper bound on lag?

Proportional Share Resource Allocation
Timing analysis



- Is a task guaranteed to complete before its deadline?
  - How late can a task be?
- Theorem: *Let c be the size of the current request of task T. Task T's lag is bounded by*

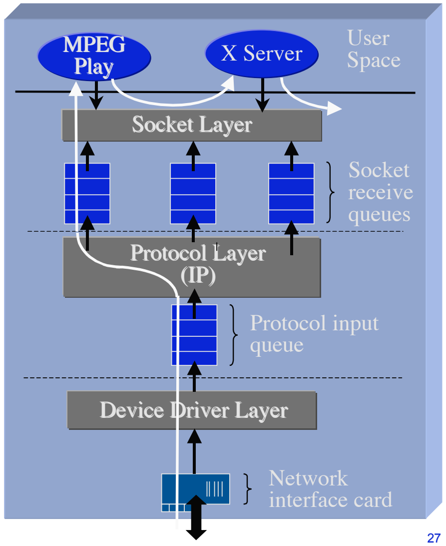
$$-q < lag_T(t) < q$$

Rate-Based Resource Allocation
FreeBSD implementation

- We've implemented RBE and proportional share scheduling in FreeBSD
- Goal: Provide integrated real-time computation and communication services in a time-shared operating system
- Technical challenge: Scheduling OS services

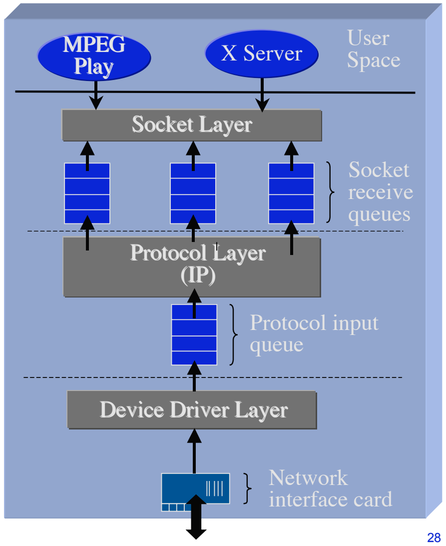
Rate-Based Resource Allocation
Integrated real-time resource allocation example

- Data arrives for a video conference over the network
- It is processed by the operating system and delivered to the application
- The application further processes and sends to the window system
- The window system paints the screen


MPEG Play
X Server
User Space
Socket Layer
Socket receive queues
Protocol Layer (IP)
Protocol input queue
Device Driver Layer
Network interface card

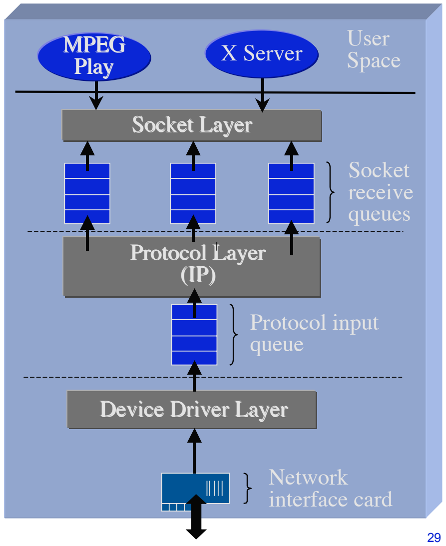
Rate-Based Resource Allocation
Integrated real-time resource allocation example

- Technical challenges:
  - Device scheduling and protocol processing
  - Application and system call scheduling
- Candidate technologies
  - Proportional share scheduling (EEVDF)
  - Constant Bandwidth Servers (CBS)
  - Rate-Based extensions to Liu and Layland (RBE)


MPEG Play
X Server
User Space
Socket Layer
Socket receive queues
Protocol Layer (IP)
Protocol input queue
Device Driver Layer
Network interface card

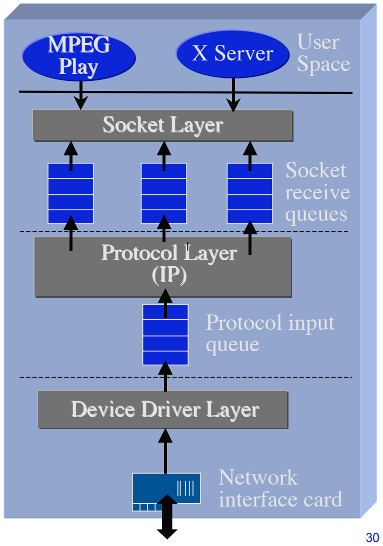
Rate-Based Resource Allocation
Integrated real-time resource allocation example

- Our study:
  - Compare the performance of applications of rate-based scheduling technology at various levels in the kernel
  - For various characterizations of real-time processing work-loads
    - » Well-behaved periodic job/task arrivals
    - » Bursty job/task arrivals
    - » “Misbehaved” job/task arrivals



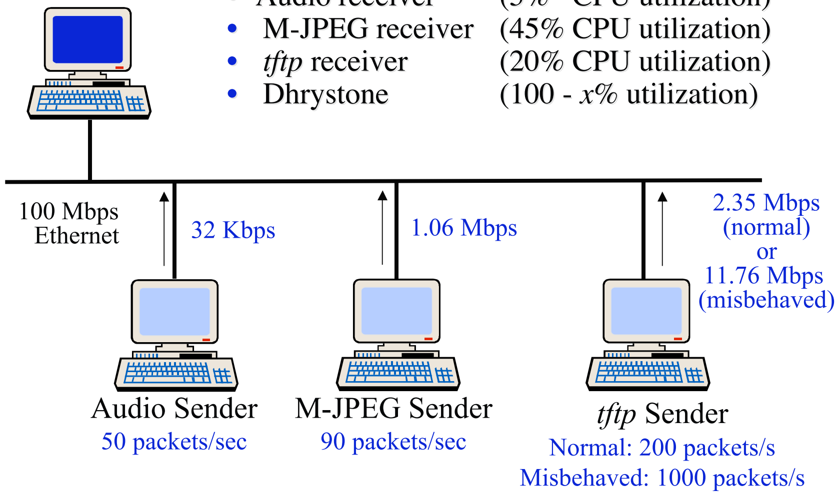
Empirical Comparisons
Experimental setup

- Modify FreeBSD UNIX to support rate-based scheduling in the “top” and “bottom” halves of the kernel
- Consider the performance of each rate-based scheme in isolation and in combinations
  - Consider the performance across a variety of multimedia workloads



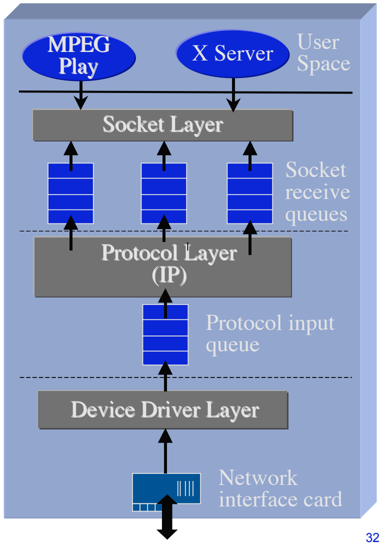
Experimental Setup
Workload generation

- Audio receiver (5% CPU utilization)
- M-JPEG receiver (45% CPU utilization)
- *tftp* receiver (20% CPU utilization)
- Dhrystone (100 - *x*% utilization)



Empirical Comparisons
Performance metrics setup

- Packets dropped at the IP layer
- Packets dropped at the socket layer
- Packets delivered to the application
- Dhrystone performance
- NIC to application response time
- Deadline miss percentage

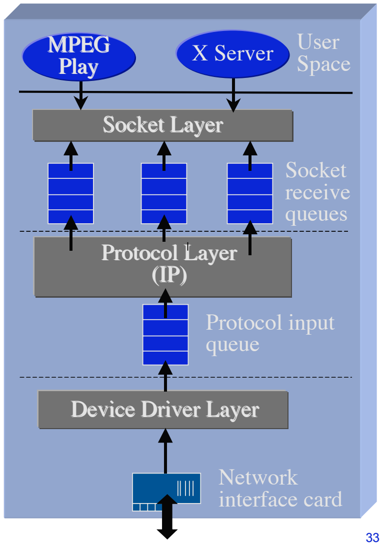


Empirical Comparisons
Experimental plan

- First consider using only
  - Proportional share,
  - CBS, and
  - RBE

  scheduling for all resource allocation problems

- Then attempt to match algorithms to the specific allocation problems where they are best suited



Experimental Results Summary
Well-behaved, periodic packet arrivals

	Prop Share			**CBS**			**RBE**		
Phone	0	0	2,993	0	0	2,977	0	0	3,000
ftp	0	0	11,961	2	0	11,914	0	0	11,944
M-JPEG	0	0	5,346	0	0	5,388	0	0	5,443

IP Drops Socket Drops Packets Delivered

- In isolation, all rate-based schemes give "perfect" (or very good) performance
  - No packets are dropped
- Liu & Layland rate-based scheduling (RBE) provides the best response times
  - (Not surprising)

Experimental Results Summary
Bursty (pareto) packet arrivals

	Prop Share			**CBS**			**RBE**		
Phone	1,585	0	1,312	0	0	2,938	0	0	3,027
ftp	5,315	0	5,408	5	0	10,760	0	0	10,778
M-JPEG	2,705	0	2,498	0	0	3,192	0	0	5,287

IP Drops Socket Drops Packets Delivered

- Proportional share scheduling degrades the performance of all applications uniformly
  - A (bad) artifact of quantum-based allocation
- CBS and RBE smooth the arrival process
  - Event driven scheduling works well here
  - Pure event-driven scheduling (RBE) gives lowest response times

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Experimental Results Summary
"Misbehaved" *ftp* packet arrivals

	Prop Share			**CBS**			**RBE**		
Phone	5	0	2,997	0	0	2,978	0	0	2,998
ftp	17,999	0	11,902	17,880	0	12,120	0	9,052	20,794
M-JPEG	56	0	5,390	0	0	5,391	0	0	5,444

IP Drops Socket Drops Packets Delivered

- Proportional share and CBS provide excellent protection/isolation for well-behaved tasks
  - *ftp* packets dropped at the IP layer
- RBE scheduling drops *ftp* packets at the socket layer
  - Pure event-driven scheduling provides no isolation
  - Dhrystone performance suffers drastically

Initial Experiments Summary
So what?

- When workload is well-behaved all schemes perform well
- Pure-event driven scheduling and quantum allocation don’t work well for “bottom-half” kernel processing
- Server-based allocation doesn’t work well for application-level processing

Combine the scheduling schemes to better match the processing requirements at each level in the system

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Combining Allocation Policies
Getting the best of all worlds

- CBS+Proportional Share scheduling

	Constant Rate			Bursty			Misbehaved		
Phone	0	0	2,869	0	0	2,998	0	0	2,797
ftp	0	0	11,722	0	0	10,340	17,898	0	11,545
M-JPEG	0	0	5,343	0	0	4,951	0	0	5,398

- RBE+Proportional Share scheduling

	Constant Rate			Bursty			Misbehaved		
Phone	0	0	2,873	0	0	2,954	0	0	2,789
ftp	0	0	11,802	0	0	10,437	17,872	0	11,647
M-JPEG	0	0	5,324	0	0	4,956	0	0	5,393

IP Drops

Socket Drops

Packets Delivered

Rate-Based Resource Allocation
Conclusions

- “One size does not fit all” (unless the external environment is (perfectly) well-behaved)
  - Quantum allocation within the kernel leads to coarse-grained control
  - Server-based allocation impractical for applications
  - Pure event scheduling doesn’t provide isolation
- Different scheduling algorithms work best at different levels of the kernel
  - Event scheduling best at the device layer
  - Server/quantum scheduling best at the application/ system call layer

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Rate-Based Resource Allocation
Summary

- There’s life beyond rate monotonic scheduling
- Rate-based resource allocation simplifies systems wherein
  - Work is generated at non-periodic but structured rates
  - Tasks may “misbehave”
- Liu and Layland extensions
  - Rate models demonstrate a fundamental distinction between static priority and deadline scheduling methods
- Fluid flow models
  - Real-time ±*quantum*
  - No fundamental distinction between real-time and non-real-time tasks
  - Provide strict isolation between tasks