Thread Synchronization: Too Much Milk

Implementing Critical Sections in Software Hard

- The following example will demonstrate the difficulty of providing mutual exclusion with memory reads and writes
  - Hardware support is needed
- The code must work *all* of the time
  - Most concurrency bugs generate correct results for *some* interleavings
- Designing mutual exclusion in software shows you how to think about concurrent updates
  - Always look for what you are checking and what you are updating
  - A meddlesome thread can execute between the check and the update, the dreaded race condition

Thread Coordination

Too much milk!

Jack
- Look in the fridge; out of milk
- Go to store
- Buy milk
- Arrive home; put milk away

Jill
- Look in fridge; out of milk
- Go to store
- Buy milk
- Arrive home; put milk away
- Oh, no!

Fridge and milk are shared data structures

Formalizing "Too Much Milk"

- Shared variables
  - "Look in the fridge for milk" – check a variable
  - "Put milk away" – update a variable
- Safety property
  - At most one person buys milk
- Liveness
  - Someone buys milk when needed
- How can we solve this problem?

How to think about synchronization code

- Every thread has the same pattern
  - Entry section: code to attempt entry to critical section
  - Critical section: code that requires isolation (e.g., with mutual exclusion)
  - Exit section: cleanup code after execution of critical region
  - Non-critical section: everything else
- There can be multiple critical regions in a program
  - Only critical regions that access the same resource (e.g., data structure) need to synchronize with each other

```
while(1) {
    Entry section
    Critical section
    Exit section
    Non-critical section
}
```


The correctness conditions

- Safety
  - Only one thread in the critical region
- Liveness
  - Some thread that enters the entry section eventually enters the critical region
  - Even if some thread takes forever in non-critical region
- Bounded waiting
  - A thread that enters the entry section enters the critical section within some bounded number of operations.
- Failure atomicity
  - It is OK for a thread to die in the critical region
  - Many techniques do not provide failure atomicity

```
while(1) {
    Entry section
    Critical section
    Exit section
    Non-critical section
}
```

Too Much Milk: Solution #0

```
while(1) {
  if (noMilk) {        // check milk (Entry section)
    if (noNote) {      // check if roommate is getting milk
      leave Note;      //Critical section
      buy milk;
      remove Note; // Exit section
    }
   // Non-critical region
}
```

- Is this solution
  - 1. Correct
  - 2. Not safe
  - 3. Not live
  - 4. No bounded wait
  - 5. Not safe and not live

What if we switch the order of checks?

- It works sometime and doesn't some other times
  - Threads can be context switched between checking and leaving note
  - Live, note left will be removed
  - Bounded wait ('buy milk' takes a finite number of steps)

Too Much Milk: Solution #1

```
turn := Jill  // Initialization
```

```
while(1) {
  while(turn ≠ Jack) ; //spin
  while (Milk) ; //spin
  buy milk;     // Critical section
  turn := Jill  // Exit section
  // Non-critical section
}
```

```
while(1) {
  while(turn ≠ Jill) ; //spin
  while (Milk) ; //spin
  buy milk;
  turn := Jack
  // Non-critical section
}
```

- Is this solution
  - 1. Correct
  - 2. Not safe
  - 3. Not live
  - 4. No bounded wait
  - 5. Not safe and not live
- At least it is safe

Solution #2 (a.k.a. Peterson's algorithm): combine ideas of 0 and 1

Variables:

- in_i: thread T_i is executing , or attempting to execute, in CS
- *turn*: id of thread allowed to enter CS if multiple want to

Claim: We can achieve mutual exclusion if the following invariant holds before thread i enters the critical section:

$$\{(\neg in_j \vee (in_j \wedge turn = i)) \wedge in_i\}$$

$$((\neg in_0 \vee (in_0 \wedge turn = 1)) \wedge in_1) \wedge ((\neg in_1 \vee (in_1 \wedge turn = 0)) \wedge in_0) \Rightarrow ((turn = 0) \wedge (turn = 1)) = \text{false}$$

Intuitively: j doesn't want to execute or it is i's turn to execute

Peterson's Algorithm

$in_0 = in_1 = \text{false};$

```
Jack
while (1) {
   in0:= true;
   turn := Jill;
   while (turn == Jill
       && in1) ;//wait
   Critical section
   in0 := false;
   Non-critical section
}
```

```
Jill
while (1) {
   in1:= true;
   turn := Jack;
   while (turn == Jack
       && in0);//wait
   Critical section
   in1 := false;
   Non-critical section
}
```

Spin!

turn=Jack, in_0 = false, in_1:= true

Safe, live, and bounded waiting
But, only 2 participants

Too Much Milk: Lessons

- Peterson's works, but it is really unsatisfactory
  - Limited to two threads
  - Solution is complicated; proving correctness is tricky even for the simple example
  - While thread is waiting, it is consuming CPU time
- How can we do better?
  - Use hardware to make synchronization faster
  - Define higher-level programming abstractions to simplify concurrent programming