COMP 633 - Parallel Computing

Lecture 9 September 16, 2021

SMM (4)

OpenMP Case Study: The Barnes-Hut N-body Algorithm

Topics

- Case study: the Barnes-Hut algorithm
 - Study an important method in scientific computing
 - » efficient n-body simulation with long range forces
 - Investigate parallelization and implementation in a shared memory multiprocessor
 - » expression and management of parallelism
 - » memory hierarchy tuning

N-body simulations: self-gravitating systems







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The *n*-body simulation problem

Simulate the evolution of a system of n bodies over time

- Pairwise interaction of bodies
 - » force f(i,j) on body i due to body j
 - » total force f(i) on body i due to all bodies
 - » acceleration of body *i* via f = ma
- Numerical integration of body velocities and positions
 - » timestep Δt

Non-negligible long-range forces

- for uniformly distributed bodies in 3D, total force due to all bodies at a given distance *r* is constant
 - cannot ignore contribution of distant bodies

• Examples

- astrophysics (gravity)
- molecular dynamics (electrostatics)

Ex: Gravitation $r_{ij} = \| \mathbf{p}_i - \mathbf{p}_j \|$ $f(i, j) = -G \cdot \frac{m_i \cdot m_j}{r_{ij}^2} \cdot \frac{\mathbf{p}_i - \mathbf{p}_j}{r_{ij}}$ $f(i) = \sum_{j \neq i} f(i, j)$

the basic simulation algorithm:

Direct approach:

 $O(n^2)$ interactions per time-step

Reducing the number of interactions

Exploit combined effect of "distant" bodies



apply this idea recursively:

- determines control-structure
- requires hierarchical decomposition of space

Formally

• *Monopole approximation* of the force on the earth due to interaction with all masses in the *Andromeda* galaxy

$$f(b_{\text{earth}}) \approx -G \frac{m_{\text{earth}} M(\mathbf{p}_{\text{earth}} - \mathbf{c})}{r^3}$$

- Monopole approximation saves work if it can be reused with multiple bodies
- Accuracy of approximation improves with
 - increasing r
 - decreasing d
 - order of the approximation
 - » Monopole, dipole, quadropole, ...
 - uniformity of body distribution



Hierarchical decomposition of space



an octree decomposition



3D

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The Barnes-Hut algorithm

stepSystem():

// P(i) is coordinates and mass of body i T := makeTree(P(1:n))forall $1 \le i \le n$ do f(i) = gravCalc(P(i), T) $\langle update \ velocities \ and \ positions \rangle$

function gravCalc(body p, treenode q)
if ("q is a leaf") then

 $\langle return \ body-body \ interaction \ (p,q) \rangle$

el se

```
\begin{array}{l} \textbf{if} (``p \ \textit{is distant enough from q"}) \ \textbf{then} \\ \langle \textit{return body-cell interaction (p,q)} \rangle \end{array}
```

el se

forall q' ∈ nonemptyChildren(q) do
 accumulate gravCalc(p,q')
 ⟨return accumulated interaction⟩

end if

end if

interaction in the case of gravitation:

$$F = G \cdot \frac{m_p \cdot m_q}{r_{pq}^2} \cdot \left[\frac{x_p - x_q}{r_{pq}}, \frac{y_p - y_q}{r_{pq}}, \frac{z_p - z_q}{r_{pq}} \right]$$

$$r_{pq} = \sqrt{(x_p - x_q)^2 + (y_p - y_q)^2 + (z_p - z_q)^2}$$

body-body interaction: use masses of bodies and distance between them.

body-cell interaction: use mass of body and mass of cell and distance between body and center of mass of cell.

force is additive; individual contributions can be accumulated.

The Barnes-Hut algorithm - Performance issues

stepSystem(P(1:n))

-- P(1:n) is sequence of bodies T := makeTree(P(1:n))forall $1 \le i \le n$ do f(i) := gravCalc(P(i), T)

 $\langle update \ velocities \ and \ positions \rangle$

function gravCalc(p, q)

if ("q is a leaf") then (return body-body interaction)

el se

 $\begin{array}{l} \textbf{if} (``p \ \textit{is distant enough from } q") \textbf{ then} \\ \langle \textit{return body-cell interaction} \rangle \end{array}$

el se

forall q' ∈ nonemptyChildren(q) do
 accumulate gravCalc(p, q')
 ⟨return accumulated interaction⟩

end if

end if

Parallelism

nested parallelism

- over bodies
- over recursively divided cells

load balance

different number of interactions for different bodies

Locality

nearby bodies interact with similar set of nodes in tree



Constructing the tree

• Small fraction f of the total work

- but sequential tree construction can limit overall speedup
 - » Amdahl's law: SP < 1/f

function makeTree(P(1:n))
for i := 1 to n do
T := insert(P(i),T)

```
\langle compute monopole approximation at each node \rangle
```

- Computing monopole approximation for each cell
 - Post-order traversal of tree
 - » At leaves, monopole coincides with single body
 - » At interior nodes, monopole is weighted sum of all children's monopoles

```
< determine child S of T in which p belongs >
```

- S' := insert(p, S)
- \langle return T with S replaced by S' \rangle

endi f



• when is a cell "distant enough"?



original criterion used by Barnes-Hut:

$$\frac{d}{r} < \theta \equiv r > \frac{d}{\theta}$$

where usually

 $0.7 \le \theta \le 1.0$

• problem: detonating galaxy anomaly



(one) solution: add distance between center of mass (cm) and geometric center of cell (c)

$$r > \frac{d}{\theta} + |cm - c|$$

Effects of acceptance criterion ... on runtime



FIG. 3.—Scaling of CRAY X-MP CPU time (CPU seconds per step per particle) for spherical, isotropic Plummer models, as a function of the number of particles, for values of the clumping parameter θ in the range $0 \le \theta \le 1.5$. Only monopole terms have been included in the force computation.

Source: L. Hernquist. Performance characteristics of tree codes. Astrophysical Journal Supplement Series, Vol. 64, Pages 715-734, 1987.

Effects of acceptance criterion ... on accuracy



FIG. 6.—Magnitude of the typical error (in percent) in the tree force computation, relative to a direct sum, as a function of θ , for selected values of the particle number N. The calculations have assumed spherical, isotropic Plummer models with softening parameter $\varepsilon = 0$, and only monopole terms have been included in the force computations.

Source: L. Hernquist. Performance characteristics of tree codes. Astrophysical Journal Supplement Series, Vol. 64, Pages 715-734, 1987.

1% accuracy sufficient for most astrophysical simulations. Different techniques with better error control necessary for other systems (*fast multipole methods*).

Effect of body distribution ... on total work



For fixed *n*

- uniform distributions generate high interaction work (shallow trees)
- non-uniform distributions generate higher tree construction and lower interaction work



Complexity of Barnes-Hut

• Tree building

- cost of tree construction depends on particle distribution
 - » cost of body insertion ∞ distance to root
 - » for a uniform distribution of *n* particles, sequential construction of the tree is $O(n \log n)$ time
- In a simulation, tree could be maintained rather than reconstructed each time step

• Force calculation (uniform distribution of bodies in 2D)

- consider computing the force acting on a body in the lower right corner
- if $\theta = 1.0$ the 3 undivided top-level squares will satisfy the acceptance criterion
- The remaining square does not satisfy the criterion, hence we descend into the next level
- each level of the tree incurs a constant amount of work while descending along the path to the lower right corner
- for a uniform distribution of *n* bodies, the length of the path is $O(\log_4 n)$
- computing the forces on n bodies is $O(n \log n)$ work
- non-uniform distribution more difficult to analyze

• Accuracy and complexity are difficult to control



Implementation issues - parallelization

• parallelization of the force computation loop:



• observations:

- force computation scales reasonably up to 16 processors
- dynamic scheduling important
- single processor performance not impressive

Processors

Implementation issues - tuning of gravCalc (1)

performance analysis of gravCalc shows

- poor cache reuse (90% L1 and 88% L2)
- poor use of floating point units
- poor reuse of subexpressions

compiler can't generate good code?

manual tuning of gravCalc

- inline computation of acceptance criterion
- inline computation of interaction
- reuse distance vector (body-cell)
- fuse loops

significant performance improvement!

- observations:
 - 2.5 times faster
 - good scaling
 - better use of FPUs and better prediction

cache reuse (93% L1 and 94% L2) still bad

```
RECURSIVE SUBROUTINE gravCalc(p,q)

IF ("q is a body") THEN

⟨compute body-body interaction; accumulate⟩

ELSE

IF ("p is distant enough from q") THEN

⟨compute body-cell interaction; accumulate⟩

ELSE

D0 q' ∈ nonemptyChildren(q)

CALL gravCalc(p,q')

END D0

END IF

END IF

END SUBROUTINE gravCalc
```

Results on O2000 (evans) for 1M particles



Processors

Implementation issues - tuning of gravCalc (2a)

• how can we improve cache reuse?

 neighboring bodies in space will most likely interact with the same cells and bodies!

sort bodies according to some spatial order:

- precompute spatial order such as Morton order or Peano-Hilbert order
- or simply order bodies as they are encountered during a depth-first treewalk of T
- Sorted bodies may also speed up subsequent tree rebuilding



Morton order





Tree order



Peano-Hilbert order

Implementation issues - tuning of gravCalc (2b)

observations:

- 30-40% increase in performance
- very good scaling
- L2 reuse now up at 99.8%
- L1 still at 93%

stepSystem(P(1:n))

T := makeTree(P(1:n)) re-order P(1:n) according to T $forall 1 \le i \le n do$ f(i) := gravCalc(P(i), T) $\langle update \ velocities \ and \ positions \rangle$

Results on O2000 (evans) for 1M particles



Processors



Implementation issues - tuning of gravCalc (3)

How can we improve L1 reuse?

- interact a group of bodies with a cell or body!
- walk the tree and compute forces for a set of neighboring bodies

```
RECURSIVE SUBROUTINE gravCalc(set P, node q)
```

```
IF ("q is a body") THEN
```

```
\textbf{D0} \hspace{0.1in} p \hspace{0.1in} \in \hspace{0.1in} P
```

 $\langle compute body-body interaction; accumulate \rangle$

```
END DO
ELSE
```

```
P' = \emptyset
```

```
DO \mathbf{p} \in \mathbf{P}
```

IF ("p is distant enough from q") THEN (compute body-cell interaction; accumulate)

ELSE

```
\mathbf{P'} = \mathbf{P'} \cup \{\mathbf{p}\}
```

```
END IF
```

```
END DO
```

```
IF (P'.NE. \emptyset) THEN
DO q' \in nonemptyChildren(q)
```

```
CALL gravCal c(\mathbf{P}', \mathbf{q}')
```

```
END DO
```

```
END IF
```

```
END IF
```

```
END SUBROUTINE gravCalc
```

Results on O2000 (evans) for 1M particles



Processors

observations:

- 20-40% increase in performance
- L1 reuse now at 99.7% (32 bodies per group)
- L2 down slightly at 96%
- ordered particles essential



Implementation issues - tuning of gravCalc (4)

Another technique to improve L1 reuse

- allow leaf-cells to contain more than 1 body
- compute the body-body interactions in a doubly nested loop.

```
RECURSIVE SUBROUTINE gravCalc(set P, node q)
  \mathbf{P}' = \emptyset
  DO p \in P
     IF ("p is distant enough from q") THEN
        (compute body-cell interaction; accumulate)
     ELSE
       IF ("q is a leaf") THEN
          DO \mathbf{p} \in \mathbf{P}, \mathbf{q}' \in \mathbf{q}
             (compute body-body interaction; accumulate)
          END DO
        ELSE
          P' = P' \cup \{p\}
        END IF
     END IF
  END DO
  IF (P'. NE. \varnothing) THEN
     D0 q' \in nonemptyChildren(q)
        CALL gravCal c(P', q')
     END DO
  END IF
END SUBROUTINE gravCalc
```

Results on O2000 (evans) for 1M particles



Processors

observations:

10% increase in performance

this algorithm will perform strictly more work than the previous versions! More particles per leaf potentially causes more body-body interactions and fewer bodycell interactions to be computed.



Implementation issues - summary

Shared memory model

- enables relatively simple parallelization of basic algorithm using OpenMP
- shared memory model critical in dynamic load balancing

Performance tuning

- overall these optimizations lead to 4-5 times faster single-processor performance
- Linear or superlinear parallel speedup to 16 processors
- optimizing serial performance is essential for obtaining good parallel performance
- last two optimization are instances of exposing parallelism to improve serial performance

Observations

- the better the performance of gravCalc the more seriously the serial tree-construction affects the overall speedup
 - » when makeTree time is included in speedup
 - speedup drops from 13.00 to 10.8 for p = 16 in first version
 - speedup drops from 15.89 to 11.74 for p = 16 on last version
- parallel tree construction algorithms!