Modeling, Identifying, and Simulating Dynamic Adaptive Streaming over HTTP (DASH)

Dynamic Adaptive Streaming over HTTP (DASH) operates by serving movie segments, each encoded at multiple bitrates, from standard web servers. During playback, successive movie segments are retrieved by a DASH application using HTTP GETs. Thus, while DASH provides industry with an effective means to serve content to users, its use presents network researchers with the following problems:

Identification

Since DASH video appears as standard HTTP traffic:

- It is hard for researchers to isolate and study DASH streams, esp if they are restricted to anonymized, header-only traces. This als means that DASH traffic has the potential to skew studies of "no browsing activity if it is not removed from a trace.
- It is difficult to apply QoS to DASH streams from unknown IP addresses without applying the same QoS to all HTTP traffic.

Question #1

Can DASH streams be distinguished from other HTTP traffic both quickly and effectively given either a live, high-speed interface or anonymized, header-only trace?

Question #2

Can the movie represented by a DASH stream be **fingerprinted** and **identified** by the **sizes** of the stream's **video segments**?



Gathering Baseline Data for the DASH Traffic Model



Andrew Reed, Jay Aikat

Department of Computer Science University of North Carolina – Chapel Hill

Research Problem

	Simulation
oecially so ormal"	Since DASH segment sizes and buffering act network conditions, captured DASH streams network simulations. Thus, to generate DASE currently presented with these three options: 1. Encode and store movies on an HTTP ser to stream and play DASH movies.
	 Labor- and storage-intensive. This appendix environment such as GENI due to stora 2. Parameterize DASH traffic. Open to criticism.
ran	 3. Stream movies from a service such as Net Expensive at scale. Introduces uncontrollable conditions to
	Question #3
ţ	Can a DASH infrastructure be replicated , and emulated , in a manner that requires little sto

* <u>Application Data Unit</u>



Preliminary Results

Overall Results (Table 1) Min Max DASH Traffic Model Measurement Average ADU Out (bytes) 433 ADU Out Standard Deviation (bytes) Average Interval (seconds) Interval Standard Deviation (seconds) 481,107 3,275,999 Max ADU In (bytes) 3,095 Average Data Rate (Kbits/s) 469 174 Data Rate Standard Deviation (Kbits/s) 2,145

Determining a Stream's Quality Level (Table 2)

	Max ADU In (B)		Avg Data Rate (Kb/s)		
Quality Level	Min	Max	Min	Max	
Best HD	2,210,954	3,275,999	2,038	3,095	
Best SD	1,642,325	1,885,911	1,620	2,089	
Better	891,322	1,086,177	622	1,551	
Good	481,107	580,330	469	1,385	

Reproducing / Identifying a Specific Movie (Fig. 1)



Future Work

Identification

Our next step is to analyze campus-wide UNC traces to better gauge the model's ability to discriminate between DASH and non-DASH traffic.

Once verified, we will incorporate the model into a program that can monitor a network link for **active DASH streams**. The monitor will be designed so that it can be easily extended by the networking community to **report** DASH streams to other devices (e.g. **OpenFlow controllers**) for routing, logging, QoS, etc.

ISP passive tap DASHmonitor Ive reporting of DASH streams						
Network Operations CenterListening Devices/Applications: • Firewalls • OpenFlow Controllers • Network Monitors						
admin@passive-tap ~ \$ adudump /dev/dag0 java DASHmonitor NEW 192.168.1.101 49621 205.128.73.253 NEW 192.168.1.143 52033 8.12.222.253 END 192.168.1.101 49621 205.128.73.253						

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DASH clients

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Instant Video.

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"Best – High Def Internet Explorer

-Outbound ADUs (HTTP GETs) Inbound ADUs (Video Segments)

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· N	Inde	l Mea	suren	nents
- 11	iouc.			
<u>g</u>	<u>Interval</u>	<u>Max</u>	Avg	<u>Data Rate</u>
rval	Std Dev	<u>ADU In</u>	<u>Data Rate</u>	Std Dev
	2	3129875	3095	1958
	2	3129875	3072	1903
	2	2927326	3071	1292
	2	2927329	3006	1419
	2	2927327	2948	1189
	3	3129875	2907	1851
	3	3233090	2826	1744
	2	3233092	2814	1713
	2	3233094	2809	1716
	3	3275999	2660	1790
	3	3275999	2655	1763
	2	2693445	2550	1729
	2	2693445	2550	1687
	2	2693445	2550	1797
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- *Table 1* shows the Min/Max for each measurement across all of the captured DASH connections.
- *Table 2* points to the possibility of identifying the quality level of a captured DASH stream.
- *Table 3* indicates that lower quality levels can be approximated by constant percentages.
- *Fig.* 1 suggests that individual movies can be reproduced/fingerprinted by their segment sizes.

Approximating Lower Quality Levels (Table 3)

Movie	Best HD	Best SD	Better	Good
7 min	102,434,048 (B)	68.3%	54.6%	32.2%
17 min	367,076,893 (B)	59.8%	36.1%	19.5%
46 min	1,055,834,938 (B)	59.5%	36.6%	20.7%
110 min	2,249,302,778 (B)	59.7%	40.8%	19.4%

Simulation

We are developing a lightweight **DASH emulator** that can replicate a **Netflix-like** architecture and generate **adaptive traffic** from any number of distributed clients streaming **real DASH movies**.

Design goals:

- The *central server* and *CDN nodes* serve all
- content as static files from standard web servers.
- Each movie is stored on the *central server* as a
- text file ordering on 7KB per hour of video. CDN nodes store a single 10-20MB dummy file
- from which the clients request byte ranges.



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