COMP 550 Algorithms and Analysis Spring 2015 Mid-Term 1 (Sample) 2015

Name _____

PID _____

Honor Pledge:

I have not given nor received unauthorized assistance in completing this exam.

Signature _____

Note:

- (1) All "lg"s are based 2 if unspecified.
- (2) Points of subproblems are evenly distributed in each problem.

- 1. (14') True or false.
 - _____(a) The sorting problem is of Ω (n)
 - ____(b) If the input list is already sorted, Quicksort runs in Θ (n lg n) time
 - ____(c) If the input list is already sorted, Randomized Quicksort runs in expectedΘ (n lg n) time
 - (d) Mergesort runs in Θ (n) best-case time.
 - (e) $f(n) = \Omega(g(n))$ implies f(n) != o(g(n)).
 - ____(f) f(n) !=0(g(n)) implies f(n) $=\Omega$ (g(n)).
 - (g) $f(n) = \Theta(g(n))$ implies f(n) != o(g(n)).
- 2. (16') For each problem, write ALL correct asymptotic relationships in {O, Θ , Ω } between the given functions. (Determine if any or some of the following relationships hold: f(x) = O(g(x)), $f(x) = \Theta(g(x))$, $f(x) = \Omega(g(x))$) (a) $f(x) = x^2-x$, $g(x) = 3x^2+20$.
 - (b) f(x) = 6x+1, $g(x) = 3 |g^2 x + 2$.
 - (c) $f(x) = 1/(\lg x), g(x) = (\lg x)/x.$
 - (d) $f(x) = x^{\lg^7} g(x) = x^3$.

You might need the following facts:

 $lg^{a}b = (log b)^{a}; lg(ab) = lga * lgb; lg_{a}b = lg_{c}b/lg_{c}a (a,b,c > 0 \& !=1)$

3. (16') Here is a recursive algorithm to solve the Tower of Hanoi problem:

Algorithm TowerofHanoi(n, Source, Spare, Destination)

Input : The n smallest rings are stacked on Source

Output: The n smallest rings are stacked on Destination

begin

TowerofHanoi(n – 1, Source, Destination, Spare)

Move the nth-largest ring directly from Source to Destination (takes time 1)

TowerofHanoi(n – 1, Spare, Source, Destination)

end algorithm

(a)What's the recurrence relation of the algorithm?

(b) How many levels are there in the recursion tree?

- (c)Obtain an asymptotic expression for T(n)?
- (d) Obtain an exact expression for T(n). _____

4. (20') Determine the *asymptotic* behavior of each of the following recurrence relations using the Master Method. Identify the values a,b, f(n), and the relationship between f(n) and n ^{lg_b(a)}. If Master method is not applicable, write "Does not apply".
(a)T(n) = 2*T(n/3) + n

(b) $T(n) = 8*T(n/2)+3n^3+4n+10$

 $(c)T(n) = 2*T(2n^{1/3})+10n$

(d) $T(n) = T(n/3) + \lg n$

- 5. (16') A "double-1" fair die when tossed will give the value 1 to 5, where the probability of 1 is twice than any other side (two sides with "1", and no "6").
 - (a) A "double-1" fair die is tossed, what is the expected value of the toss, under the condition that the value produced is odd? _____
 - (b) Suppose two "double-1" fair dice are tossed, what is the probability that they will all produce equal values?
 - (c) Suppose two "double-1" fair dice are tossed, what is the probability that the sum of their values will equal to 3? ______
 - (d) What is the expected value for the sum of two tosses of a "double-1" fair die?_____

6. (8') The following algorithm prints all divisors of the number n

```
x = 1;
while (x <= n) { // Run x = 1, 2, ..., n
if (n% x == 0) { // x is a divisor of n
System.out.println(x); // Print x (because it's a divisor)
}
x++;
}
(a) State a loop invariant for the while loop.
At the beginning of the ith loop, we have ...
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- (b) Assume the loop invariant in (a) holds, prove the correctness of the algorithm.
- 7. (10') You are given k nuts N and k bolts B, each of different diameter. For any pair (n,b) ∈ N*B, you will learn either n<b, n>b or they fit (n=b). Note that you cannot compare nut to nut or bolt to bolt directly. We'd like to find all pairs that fit.

In the case that NOT every nut has a bolt that it is paired with and vice versa, prove that any algorithm must do Ω (k²) tests in the worst case, for some input.