

A Note on the Interpretation of σ_k

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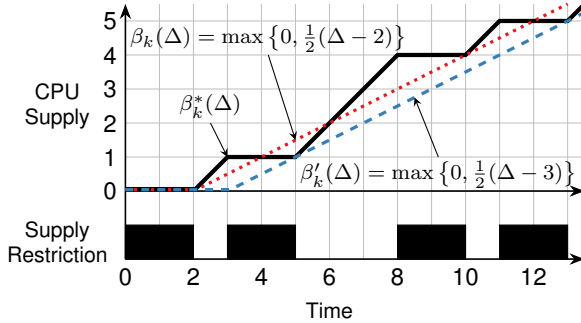


Figure 1: Service functions discussed in Example 1.

1 Problem

In several works pertaining to restricted processor supply [3, 4, 5, 6, 7, 8], a piecewise-linear lower bound called a *service function* [1] is given for the available processor time to tasks in some task set τ on processor k over any interval of length $\Delta \geq 0$. This bound is of the form

$$\beta_k(\Delta) = \max\{0, \widehat{u}_k \cdot (\Delta - \sigma_k)\},$$

and accompanied by an interpretation like the following.

\widehat{u}_k is the total long-term utilization available to the tasks in τ on processor k and σ_k is the maximum length of time when the processor can be unavailable. [3]

This interpretation of σ_k is flawed, and may result in the construction of a function $\beta_k(\Delta)$ that is not actually a lower bound to the available processor time on processor k . To see this, consider the following example.

Example 1. Figure 1 depicts a processor k with a pattern of supply restriction that repeats every 8 time units. The intervals over which the processor is unavailable are shown in black at the bottom of the figure, and the total supply is given by the curve $\beta_k^*(\Delta)$. The long-term utilization available on processor k is $1/2$. According to the interpretation above, the value for σ_k should be 2, giving the function β_k shown in the figure. Clearly, β_k is not a lower-bound of $\beta_k^*(\Delta)$ (see e.g. the interval (4, 6)).

2 Correct Explanation

A correct interpretation of the constants σ_k and \widehat{u}_k in the function $\beta_k(\Delta)$ is given by Erickson and Anderson [2], reproduced here.

σ_k is set to the x -intercept necessary in order for $\beta_k(\Delta)$ to lower-bound the actual supply, when the slope of $\beta_k(\Delta)$ is \widehat{u}_k . [2]

Example 1 (continued). Referring back to Figure 1, choosing a value of 3 for σ_k yields the function $\beta'_k(\Delta)$, which is a tight lower-bound of the actual supply $\beta_k^*(\Delta)$.

References

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