Searchable Symmetric Encryption

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Advanced Topics in Network Security Spring 2006

Yesterday

- Motivation for searchable encryption
- First SSE scheme [SWP00]
- Attacks on [SWP00]
- Conjunctive SSE [GSW04,PKL04,BKM05]

Today

- Limitations of Song et al.'s security model
- More formal work on SSE [Goh03,CM05]
- New definitions

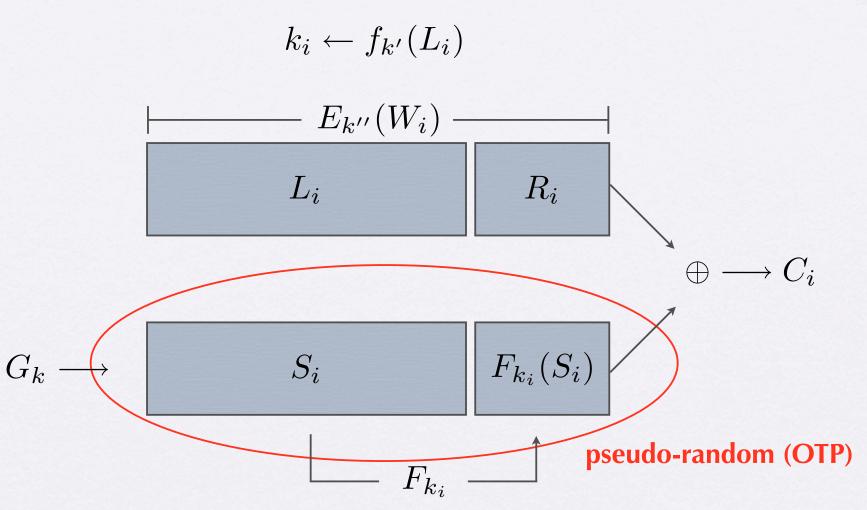
Practical Techniques[swp00]

- Song et al. provide proofs of security
 - "Our techniques are provably secure" (p. 1)
- Yet
 - there are statistical attacks
 - leaks location of words

What's Going on?

- Are the proofs wrong?
- What are they proving?
- Is it meaningful?

What are they Proving?



- Is proving that the key stream is pseudo-random useful?
- Depends on the adversarial model!

Adversarial Model

Who are we protecting against?

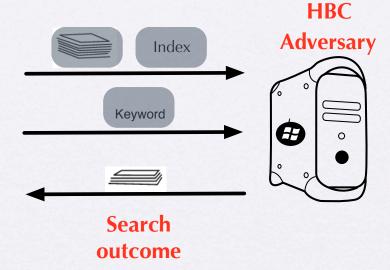
the server

What are its goal?

info. about documents and keywords

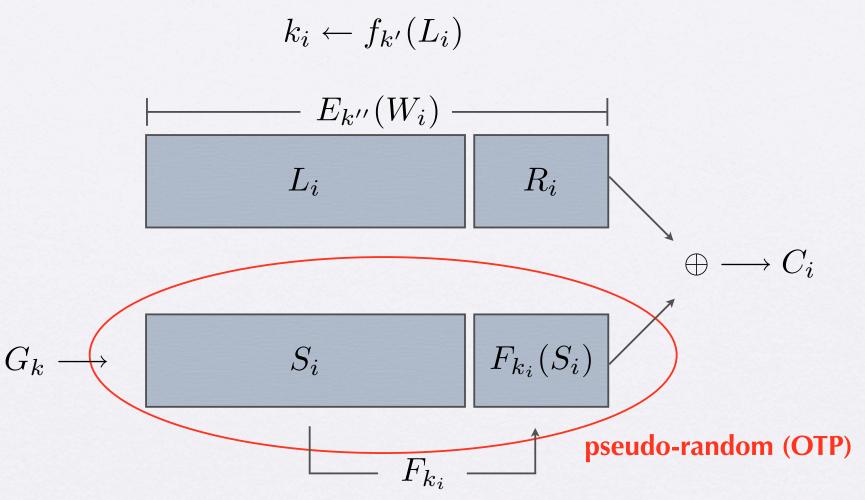
How much power does it have?





it can search!

What are they Proving?



	Ideal model	[SWP00]	
Adversary	server	server	
Adv.'s Goal	recovering documents & keywords	recovering documents & keywords	
Adv.'s Power	it can search	none	
Meaning	documents and keywords are secure against server that can search	documents are secure against server that cannot search	

Secure Indexes [Goh03]

- Introduces a stronger (better) security model
 - IND2-CKA: security against chosen-keyword attacks
- Provides provably secure and efficient construction
 - separates index from ciphertext
 - one index per document
 - based on pseudo-random functions & Bloom filters

Adversarial Model

Who are we protecting against?

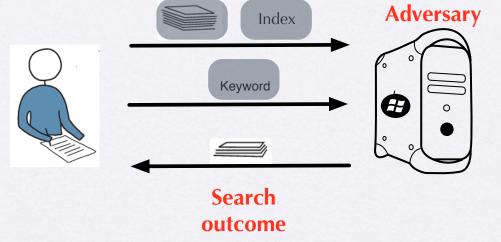
the server

What are its goals?

info. about documents and keywords

How much power does it have?

it can search!



Formalizing the Adversarial Model

 How exactly do we capture the adversarial model formally?

Adversarial Model

Who are we protecting against?

the server

What are its goals?

info. about documents and keywords

How much power does it have?

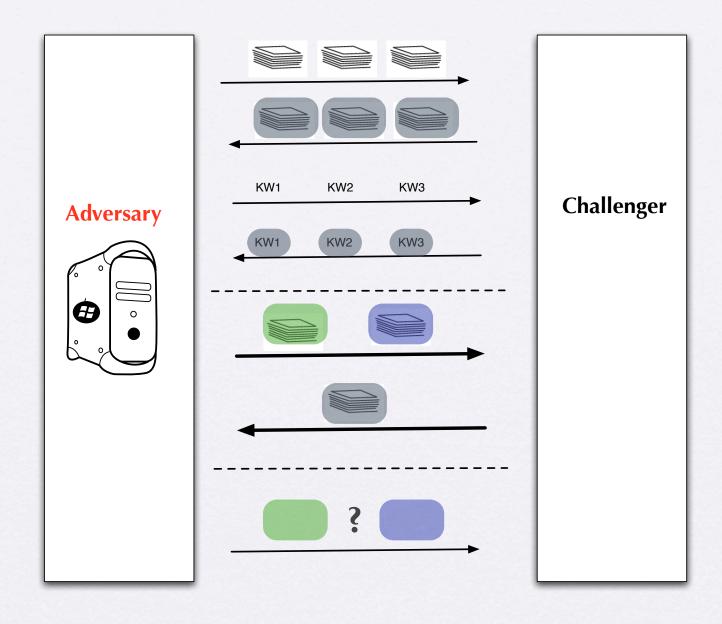
it can search!

Probabilistic polynomial-time (PPT) algorithm

indistinguishability

allow adversary to generate and search many documents and keywords

IND2-CKA



	Ideal model	[SWP00]	IND2-CKA
Adversary	server	server	server
Adv.'s Goal	recovering recovering documents & documents & keywords		recovering documents
Adv.'s Power	it can search	none	it can search
Meaning	documents and keywords are secure against server that can search	documents are secure against server that cannot search	documents are secure against server that can search

Secure Indexes [Goh03]

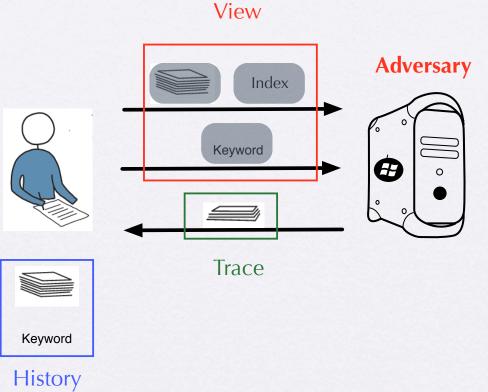
- Limitations:
 - IND2-CKA says nothing about trapdoors
 - One has to prove IND2-CKA + security of trapdoors

Privacy Preserving [CM05]

- Introduces a stronger security model than IND2-CKA
 - CM: security of index and trapdoors against chosenkeyword attacks
- Provides provably secure constructions
 - separates index from ciphertext
 - one index per document
 - Pseudo-random functions

CM Security [CM05]

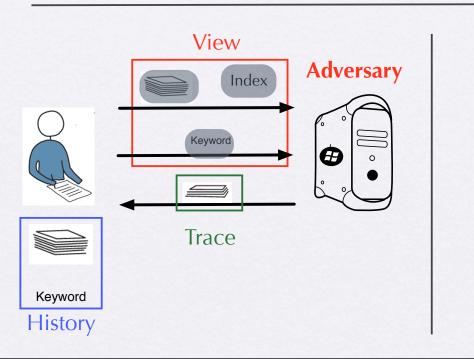
- **History**: documents and words queried
- View: what the server sees
- **Trace**: minimum information leaked

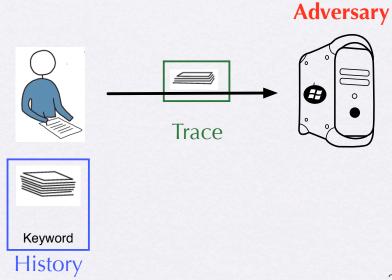


CM Security [CM05]

• for all q, for all adversaries, for any function f, there exists a simulator such that for all histories

$$\left| \Pr \left[\begin{array}{c} \mathcal{A}(\mathsf{View_q}) = \\ f(\mathsf{History_q}) \end{array} \right] - \Pr \left[\begin{array}{c} \mathcal{S}(\mathsf{Trace_q}) = \\ f(\mathsf{History_q}) \end{array} \right] \right| \leq \mathsf{negl}(k)$$





CM Security [CM05]

- Intuition: anything the adversary can recover about the history from the view, can be recovered from the trace
- Implication: no adversary can recover any information about the documents or word queries that he is not supposed to

	Ideal model	[SWP00]	IND2-CKA	CM
Adversary	server	server	server	server
Adv.'s Goal	recovering documents & keywords	recovering documents & keywords	recovering documents	recovering documents & keywords
Adv.'s Power	it can search	none	it can search	it can search
Meaning	documents and keywords are secure against server that can search	documents are secure against server that <i>cannot</i> search	documents are secure against server that <i>can search</i>	documents and keywords are secure against server that can search

- So did Chang and Mitzenmacher finally get it right?
- Not exactly...

	Ideal model	[SWP00]	IND2-CKA	CM
Adversary	server	server	server	server
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Meaning	documents and keywords are secure against server that can search	documents are secure against server that cannot search	documents are secure against server that can search	documents and keywords are secure against server that can search