



*3<sup>rd</sup>*

# Quiz 02 Review Session

COMP 210 / 2024 Summer Session I

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## Quiz 02 Format

- 30 minutes at the start of class.
- *On paper* - bring a pencil!
- Question Types:
  - Multiple choice, T/F, select all that apply, fill in the blank,
  - *No code writing on this quiz - but be able to trace given Java code!*

Possibly draw a  
diagram\* (does not have  
to be extremely  
detailed)



## Exercise Check-In Question

- ... Specifics covered in lecture (5/28)



## On Quiz 02

- Big-O Analysis
  - Analyzing code snippets for runtime analysis, including recursive code
- The `List` Abstract Data Type
  - Understand `ArrayList` and `LinkedList` on the heap
  - Explain trade-offs between both, justified using big-O notation



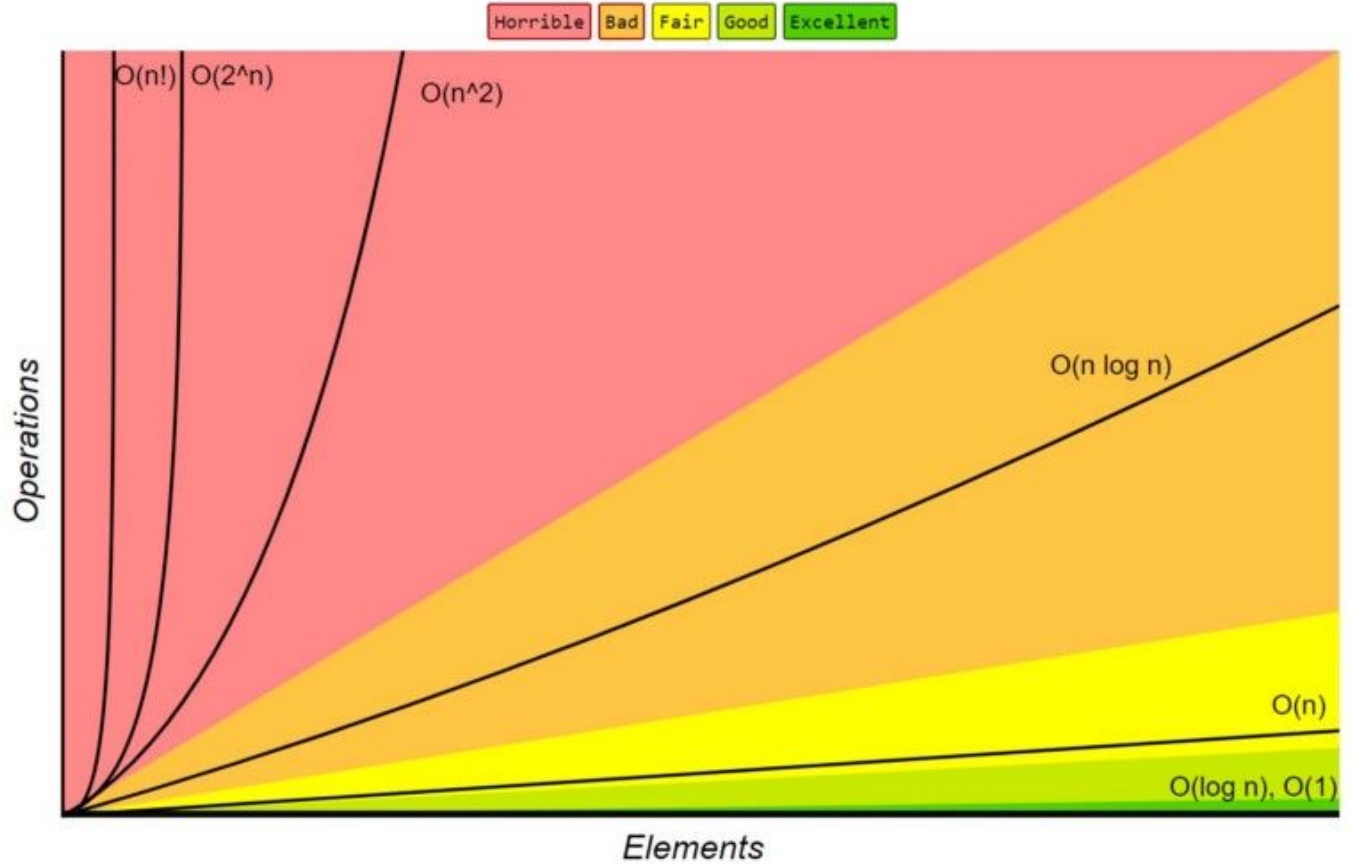
## Review: Big-O Analysis

- We need a way to determine *how efficiently* algorithms run.
  - We need notation to be able to compare the *efficiency* of algorithms.
  - This is called Big-O Notation.
- We can tell how efficient algorithms run by comparing *how many operations* an algorithm performs compared to the *number of inputs we supply to it*.

→ "time"

# Big-O Graph Comparisons

## Big-O Complexity Chart



## Recursive Example 1

```
void foo5(int n) {
```

```
    bc if(n <= 0) return 1;
```

```
    rc return 1 + foo(n-1);
```

```
}
```

$O(N) * O(1/H)$

$O(1) * O(N) = O(N)$

$C + f(n-1)$

$O(N)$

$$f(5) = 1 + f(4)$$

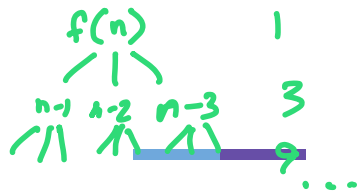
$$f(4) = 1 + f(3)$$

$$f(3) = 1 + f(2)$$

$$f(2)$$

$$f(1)$$

$$f(0)$$



$3^n \rightarrow 2^n$

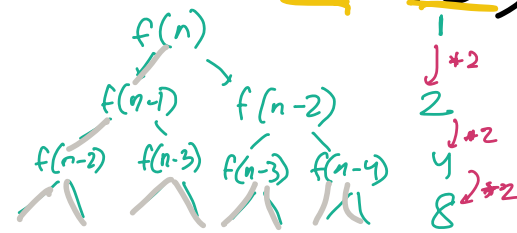
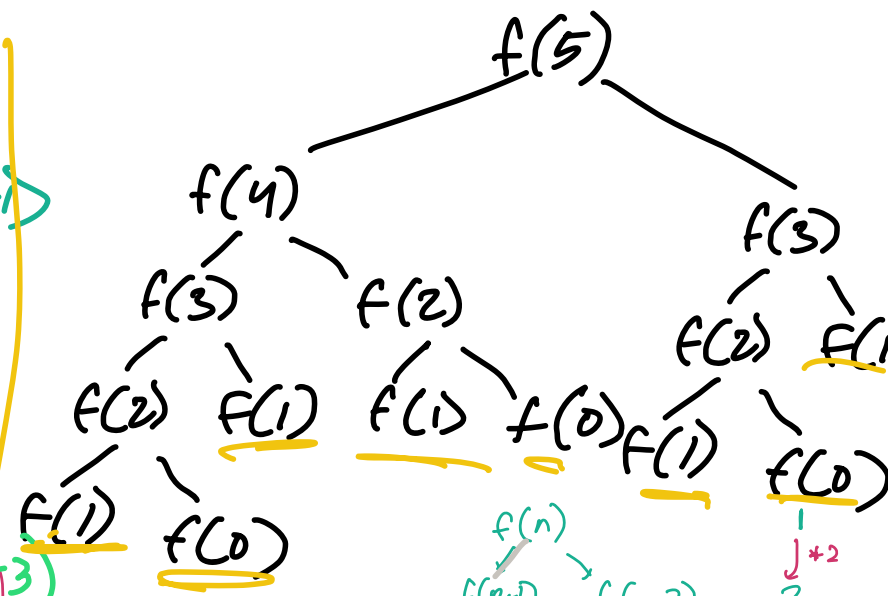
# Recursive Example 2

```

f
void fib(int n) {
    if(n < 2) return n;
    return (fib(n-1) + fib(n-2) + fib(n-3));
}

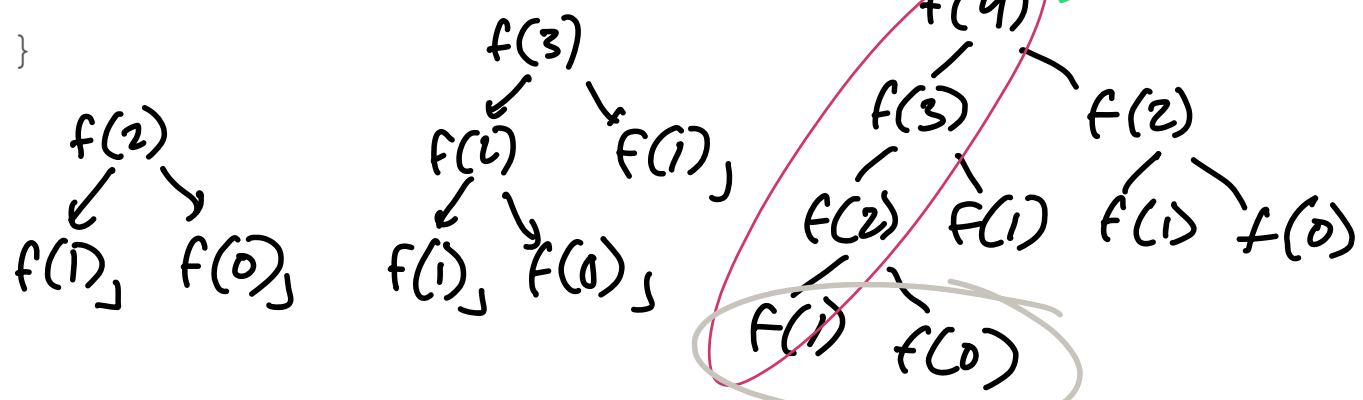
```

$c + f(n-1)$   
 $n \neq 2^n$



$2 * 2 * 2 * \dots$   
 $n$   
 $n 2^n$   
 $2^n$

$1 * 2$

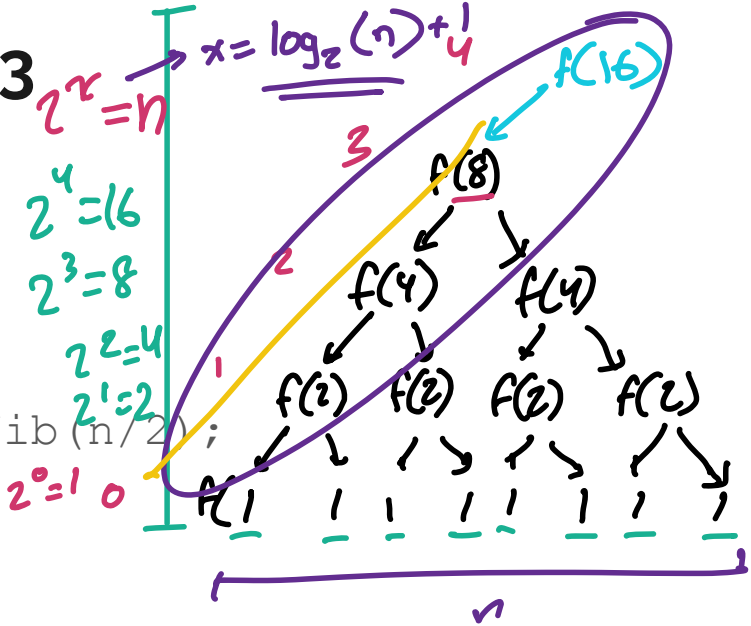




## Recursive Example 3

```
void f(int n) {  
    if (n <= 1) return n;  
    return fib(n/2) + fib(n/2);  
}
```

16



$O(n) + O(n)$

$O(n) \neq O(\log n) = \underline{O(n \log n)}$

# Recursion Big-O Guide

Given a recursive case:

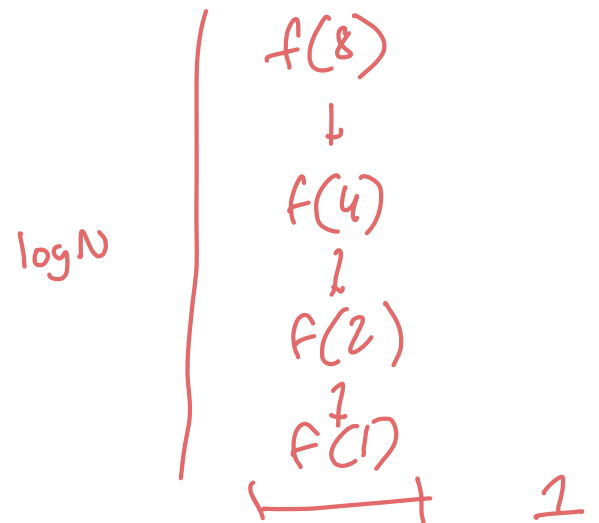
$$\textcircled{1} \underbrace{c}_{\text{constant}} \pm \underbrace{f(n-\#)}_{\text{same \#}} \xrightarrow{\text{one recursive call}} O(N)$$

Input size decreases linearly (+/-)

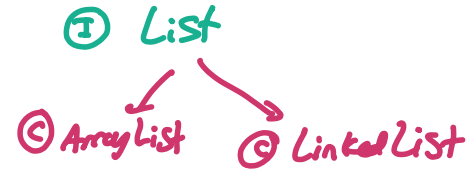
$$\textcircled{2} \underbrace{f(n-\#_1)}_{\text{any operation}} + \underbrace{f(n-\#_2)}_{\text{any operation}} + \dots \Rightarrow O(2^N)$$

$$\textcircled{3} f\left(\frac{n}{\#_1}\right) + f\left(\frac{n}{\#_2}\right) \Rightarrow O(N \log N)$$

$$\textcircled{4} c \pm f\left(\frac{n}{\#_1}\right) \Rightarrow O(\log N)$$



# ArrayList Representation

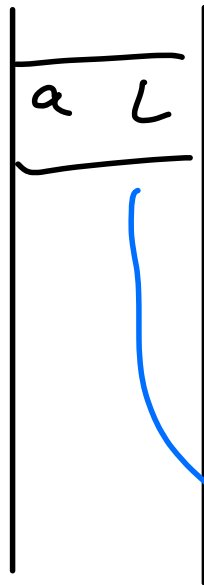


- Recall that List is an abstract data type.
- **ArrayList** is one implementation of the List interface.

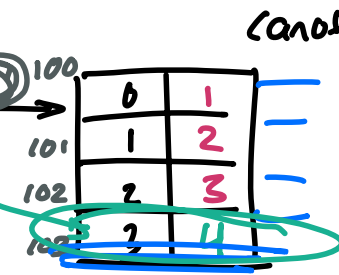
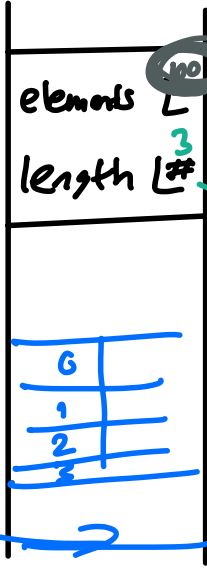
# ArrayList Representation

$a = [1, 2, 3]$

Stack

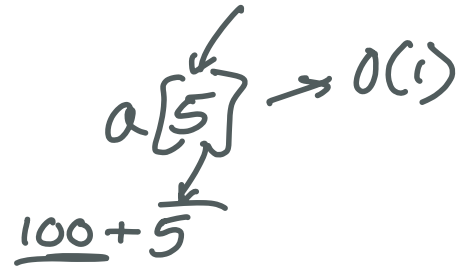


Heap



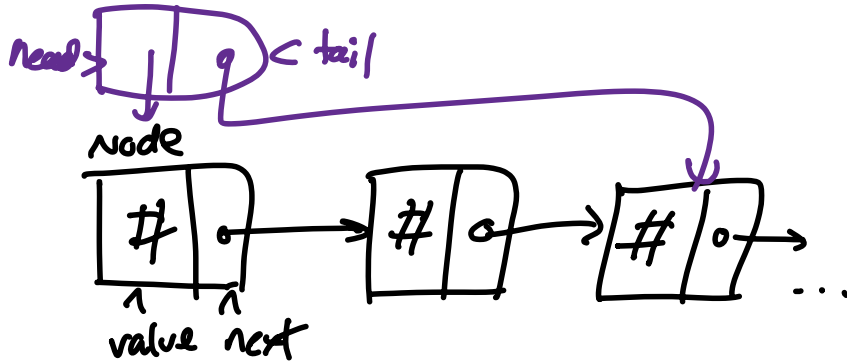
(another part of the heap)

$a.add(4)$

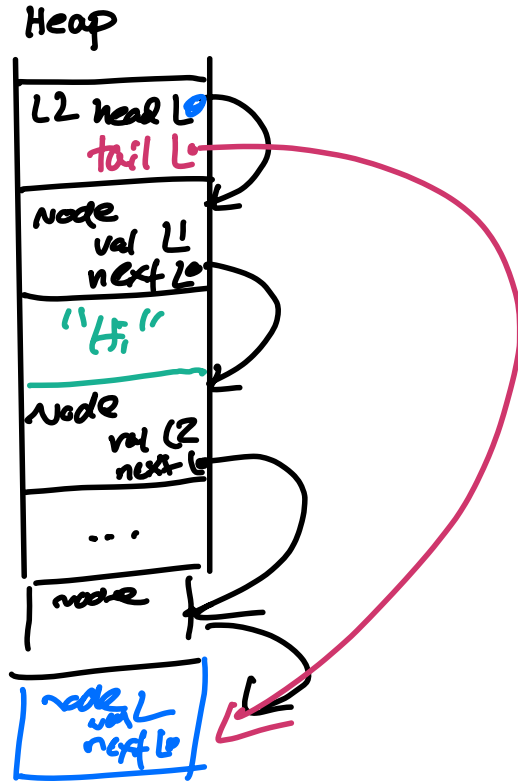


# LinkedList Representation

- `LinkedList` is another implementation of the `List` interface.



# LinkedList Representation



# Deriving List Time Complexities

\* = amortized

	get(0)	get(i)	get(n)	insert(0)	insert(i)	insert(n)	remove(0)	remove(i)
ArrayList	$O(1)$	$O(1)$	$O(1)$	$O(N)$ Avg = $O(N)$	$O(N)$ Avg = $O(N)$	$O(N)$ * Avg = $O(1)$	$O(N)$	$O(N)$
LinkedList (Head only)	$O(1)$	$O(N)$	$O(N)$ *	$O(1)$	$O(N)$	$O(n)$ *	$O(1)$	$O(N)$
LinkedList (Head and <u>Tail</u> )	$O(1)$	$O(N)$	$O(1)$ *	$O(1)$	$O(N)$	$O(1)$ *	$O(1)$	$O(N)$

$O(N)$  for  
remove(n)  
also

