

## Cheat Sheet

Dataflow Analysis:

$$\text{for}( v \in V ) \{$$
$$\quad \text{out}(v) := \bigcup_{s: s \in \text{successor}(v)} \text{in}(s)$$
$$\quad \text{in}(v) := \text{use}(v) \cup (\text{out}(v) \setminus \text{def}(v))$$
$$\}$$

Invariants:

$$I0 \equiv \text{use}(v) \subseteq \text{in}(v)$$
$$I1 \equiv \text{out}(v) \setminus \text{def}(v) \subseteq \text{in}(v)$$
$$I2 \equiv \forall s : s \in \text{successor}(v) :: \text{in}(s) \subseteq \text{out}(v)$$

Notation:

$$\text{Followers}(A) \equiv \text{FL}(A)$$
$$\text{Starters}(A) \equiv \text{ST}(A)$$
$$\text{Nullable}(A) \equiv \text{N}(A)$$

Capitals are non-terminals, lowercase are terminals.

Greek letters are sequences.

Nullable Induction:

<i>Observed</i>	<i>Rule</i>
1. $\alpha = \varepsilon$	$\text{Nullable}(\alpha) = \text{true}$
2. $\alpha = t$	$\text{Nullable}(\alpha) = \text{false}$
3. $\alpha = A$	$\text{Nullable}(\alpha) = \text{Nullable}(A)$
4. $\alpha = \alpha_1 \alpha_2 \dots \alpha_n$	$\text{Nullable}(\alpha) = \text{Nullable}(\alpha_1) \wedge \dots \wedge \text{Nullable}(\alpha_n)$
5. $\alpha = \alpha_1   \alpha_2   \dots   \alpha_n$	$\text{Nullable}(\alpha) = \text{Nullable}(\alpha_1) \vee \dots \vee \text{Nullable}(\alpha_n)$
6. $\alpha = \beta^*$	$\text{Nullable}(\alpha) = \text{true}$

## Starters Induction:

<i>Observed</i>	<i>Rule</i>
1. $\alpha = \varepsilon$	$\text{Starters}(\alpha) = \{ \varepsilon \}$
2. $\alpha = t$	$\text{Starters}(\alpha) = \{ t \}$
3. $\alpha = A$	$\text{Starters}(\alpha) = \text{Starters}(A)$
4. $\alpha = \alpha_1 \alpha_2 \dots \alpha_n$	$\text{Starters}(\alpha) = \text{Starters}(\alpha_1) \oplus \text{Starters}(\alpha_2 \dots \alpha_n)$
5. $\alpha = \alpha_1   \alpha_2   \dots   \alpha_n$	$\text{Starters}(\alpha) = \text{Starters}(\alpha_1) \cup \dots \cup \text{Starters}(\alpha_n)$
6. $\alpha = \beta^*$	$\text{Starters}(\alpha) = \text{Starters}(\beta) \cup \{ \varepsilon \}$

$$A \oplus B = \begin{cases} A & \text{if } \varepsilon \notin A \\ (A \setminus \{ \varepsilon \}) \cup B & \text{otherwise} \end{cases}$$

## Followers First Step:

$$\text{FL}_0(A) = \left( \bigcup_{C ::= \alpha A \beta} \text{ST}(\beta) \right) \setminus \{ \varepsilon \}$$

## Followers Inductive Step:

$$\text{FL}_{i+1}(A) = \text{FL}_i(A) \cup \bigcup_{C ::= \alpha A \beta \text{ and Nullable}(\beta)} \text{FL}_i(C)$$

## Followers Final Step:

$$\text{FL}(A) = \text{FL}_n(A) \cup \begin{cases} \{ \varepsilon \} & \text{if } S \Rightarrow^* \alpha A \\ \{ \} & \text{otherwise} \end{cases}$$

For each choice in  $A ::= \beta(\alpha_1 | \dots | \alpha_m)\gamma$  define:

$$\text{Predict}(\alpha_i) = \text{Starters}(\alpha_i\gamma) \oplus \text{Followers}(A)$$

For repetitions in  $A ::= \beta(\alpha)^*\gamma$

$$\text{Predict}(\alpha) = \text{Starters}(\alpha)$$

$$\text{Predict}(\gamma) = \text{Starters}(\gamma) \oplus \text{Followers}(A)$$

For sequences  $A ::= \alpha\beta\gamma$

$$\text{Predict}(A) = \text{ST}(\alpha) \oplus \text{ST}(\beta) \oplus \text{ST}(\gamma)$$

x64:

pop rm: Stores data at [rsp] into rm, and  $\text{rsp} += 8$

push rm: Does  $\text{rsp} -= 8$ , then stores rm at [rsp]

mov rm,r: Takes r, stores it in rm

mov rm,imm32: Stores imm32 in rm

RM examples:

Form	Example
r	# rdx
[rdisp + disp]	# [rbp-10]
[ridx*mult + disp]	# [rcx*8+11223344]
[rdisp + ridx*mult + disp]	# [rbx+rsi*4-8]
[disp]	# [11223344]

Useful Registers:

rax, rcx, rdx, rbx, rsp, rbp, rsi, rdi

r8, r9, r10, r11, r12, r13, r14, r15