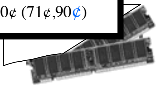
Introduction to I/O and Disk Management

Secondary Storage Management

Disks — just like memory, only different

- Why have disks?
  - Memory is small. Disks are large.
    - Short term storage for memory contents (e.g., swap space).
    - Reduce what must be kept in memory (e.g., code pages).
  - Memory is volatile. Disks are forever (?!)
    - File storage.

	GB/dollar	dollar/GB
RAM	0.013(0.015,0.01)	$77($68,$95)
Disks	3.3(1.4,1.1)	30¢ (71¢,90¢)



Capacity : 2GB vs. 1TB
2GB vs. 400GB
1GB vs 320GB

How to approach persistent storage

- Disks first, then file systems.
  - Bottom up.
  - Focus on device characteristics which dominate performance or reliability (they become focus of SW).
- Disk capacity (along with processor performance) are the crown jewels of computer engineering.
- File systems have won, but at what cost victory?
  - Ipod, iPhone, TivO, PDAs, laptops, desktops all have file systems.
  - Google is made possible by a file system.
  - File systems rock because they are:
    - Persistent.
    - Heirarchical (non-cyclical (mostly)).
    - Rich in metadata (remember cassette tapes?)
    - Indexible (hmmm, a weak point?)
- The price is complexity of implementation.

Different types of disks

- Advanced Technology Attachment (ATA)
  - Standard interface for connecting storage devices (e.g., hard drives and CD-ROM drives)
  - Referred to as IDE (Integrated Drive Electronics), ATAPI, and UDMA.
  - ATA standards only allow cable lengths in the range of 18 to 36 inches. CHEAP.
- Small Computer System Interface (SCSI)
  - Requires controller on computer and on disk.
  - Controller commands are sophisticated, allow reordering.
- USB or Firewire connections to ATA disc
  - These are new bus technologies, not new control.
- Microdrive – impressively small motors

Different types of disks

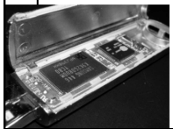
- Bandwidth ratings.
  - These are unachievable.
  - 50 MB/s is max off platters.
  - Peak rate refers to transfer from disc device's memory cache.
- SATA II (serial ATA)
  - 3 Gb/s (still only 50 MB/s off platter, so why do we care?)
  - Cables are smaller and can be longer than pATA.
- SCSI 320 MB/s
  - Enables multiple drives on same bus

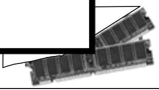
Mode	Speed
UDMA0	16.7 MB/s
UDMA1	25.0 MB/s
UDMA2	33.3 MB/s
UDMA3	44.4 MB/s
UDMA4	66.7 MB/s
UDMA5	100.0 MB/s
UDMA6	133 MB/s

Flash: An upcoming technology

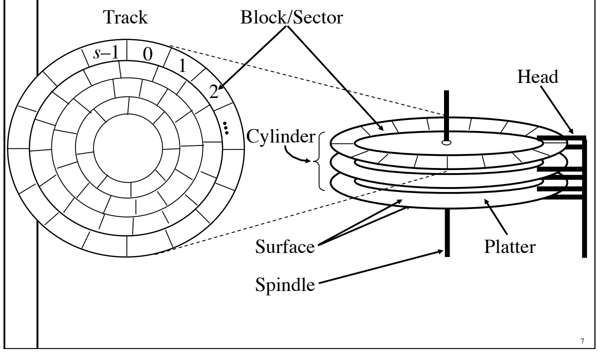
- Flash memory gaining popularity
  - One laptop per child has 1GB flash (no disk)
  - Vista supports Flash as accelerator
  - Future is hybrid flash/disk or just flash?
  - Erased a block at a time (100,000 write-erase-cycles)
  - Pages are 512 bytes or 2,048 bytes
  - Read 18MB/s, write 15MB/s
  - Lower power than (spinning) disk

	GB/dollar	dollar/GB
RAM	0.013(0.015,0.01)	$77($68,$95)
Disks	3.3 (1.4,1.1)	30¢ (71¢,90¢)
Flash	0.1	$10

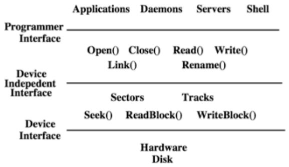




Anatomy of a Disk
Basic components



Disk structure: the big picture


Applications Daemons Servers Shell

Programmer Interface

Open() Close() Read() Write()
Link() Rename()

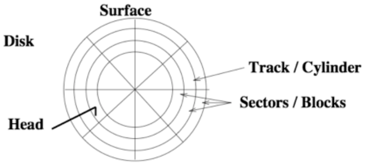
Device Indepedent Interface

Sectors Tracks
Seek() ReadBlock() WriteBlock()

Device Interface

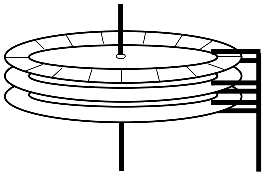
Hardware
Disk

- Physical structure of disks



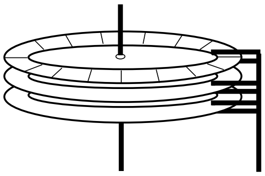
Anatomy of a Disk
Seagate 73.4 GB Fibre Channel Ultra 160 SCSI disk

- Specs:
  - 12 Platters
  - 24 Heads
  - Variable # of sectors/track
  - 10,000 RPM
    - Average latency: 2.99 *ms*
  - Seek times
    - Track-to-track: 0.6/0.9 *ms*
    - Average: 5.6/6.2 *ms*
    - Includes acceleration and settle time.
  - 160-200 MB/s peak transfer rate
    - 1-8K cache
  - 12 Arms
  - 14,100 Tracks
  - 512 bytes/sector



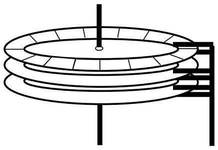
Anatomy of a Disk
Example: Seagate Cheetah ST373405LC (March 2002)

- Specs:
  - Capacity: 73GB
  - 8 surfaces per pack
  - # cylinders: 29,549
  - Total number of tracks per system: 236,394
  - Variable # of sectors/track (776 sectors/track (avg))
  - 10,000 RPM
    - average latency: 2.9 ms.
  - Seek times
    - track-to-track: 0.4 ms
    - Average/max: 5.1 ms/9.4ms
  - 50-85 MB/s peak transfer rate
    - 4MB cache
  - MTBF: 1,200,000 hours



Disk Operations
Read/Write operations

- Present disk with a sector address
  - *Old: DA* = (*drive*, *surface*, *track*, *sector*)
  - New: Logical block address (LBA)
- Heads moved to appropriate track
  - seek time
  - settle time
- The appropriate head is enabled
- Wait for the sector to appear under the head
  - “rotational latency”
- Read/write the sector
  - “transfer time”



Read time:
seek time + latency + transfer time
(5.6 *ms* + 2.99 *ms* + 0.014 *ms*)

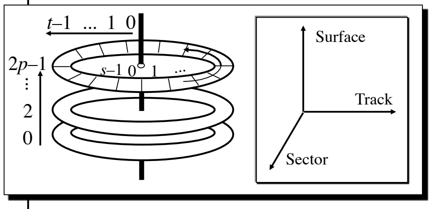
Disk access latency

- Which component of disk access time is the longest?
  - A. Rotational latency
  - B. Transfer latency
  - C. Seek latency

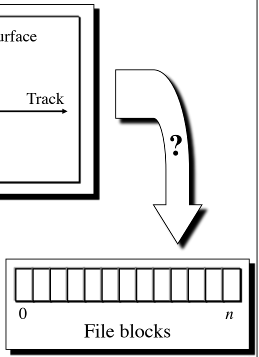
Disk Addressing

- Software wants a simple “disc virtual address space” consisting of a linear array of sectors.
  - Sectors numbered 1..N, each 512 bytes (typical size).
  - Writing 8 surfaces at a time writes a 4KB page.
- Hardware has structure:
  - Which platter?
  - Which track within the platter?
  - Which sector within the track?
- The hardware structure affects latency.
  - Reading from sectors in the same track is fast.
  - Reading from the same cylinder group is faster than seeking.

Disk Addressing
Mapping a 3-D structure to a 1-D structure



- Mapping criteria
  - block n+1 should be as “close” as possible to block n



The Impact of File Mappings
File access times: Contiguous allocation

- Array elements map to contiguous sectors on disk
  - Case1: Elements map to the middle of the disk

$$\underbrace{\underbrace{5.6}_{\boxed{\text{Seek Time}}} + \underbrace{3.0}_{\boxed{\text{Lat-ency}}}}_{\boxed{\text{Constant Terms}}} + \underbrace{\underbrace{6.0\,\frac{2{,}048}{424}}_{\boxed{\text{Transfer Time} = \left(\textit{time per revolution}\right) \times \left(\textit{number of revolutions required to transfer data}\right)}}}_{\boxed{\text{Variable Term}}} = 8.6 + 29.0 = 37.6\ ms$$

The Impact of File Mappings
File access times: Contiguous allocation

- Array elements map to contiguous sectors on disk
  - Case1: Elements map to the middle tracks of the platter

$$5.6 + 3.0 + 6.0\,\frac{2{,}048}{424} = 8.6 + 29.0 = 37.6\ ms$$

Case2: Elements map to the inner tracks of the platter

$$5.6 + 3.0 + 6.0\,\frac{2{,}048}{212} = 8.6 + 58.0 = 66.6\ ms$$

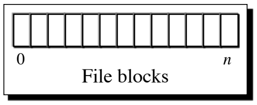
Case3: Elements map to the outer tracks of the platter

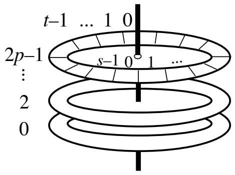
$$5.6 + 3.0 + 6.0\,\frac{2{,}048}{636} = 8.6 + 19.3 = 27.9\ ms$$

Disk Addressing
The impact of file mappings: Non-contiguous allocation

- Array elements map to random sectors on disk
  - Each sector access results in a disk seek

$$2{,}048 \times (5.6 + 3.0) = 17.6\ \textit{seconds}$$





Practical Knowledge

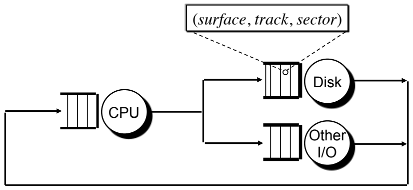
- If the video you are playing off your hard drive skips, defragment your file system.
- OS block allocation policy is complicated. Defragmentation allows the OS to revisit layout with global information.
- Unix file systems need defragmentation less than Windows file systems, because they have better allocation policies.

Defragmentation Decisions

- Files written when the disk is nearly full are more likely to be fragmented.
  - A. True
  - B. False

Disk Head Scheduling
Maximizing disk throughput

- In a multiprogramming/timesharing environment, a queue of disk I/O requests can form

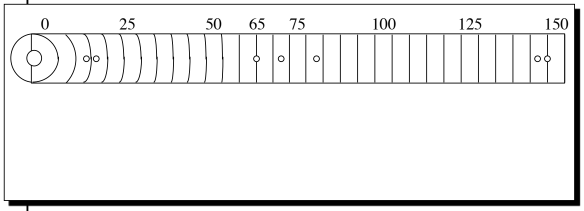


The OS maximizes disk I/O throughput by minimizing head movement through *disk head scheduling*

Disk Head Scheduling
Examples

- Assume a queue of requests exists to read/write tracks:

83 72 14 147 16 150 and the head is on track 65



Disk Head Scheduling
Examples

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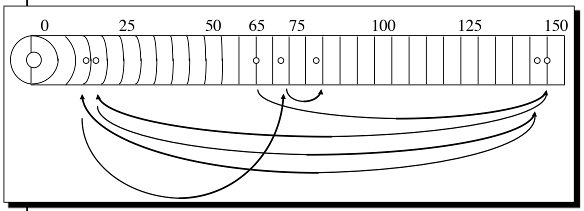


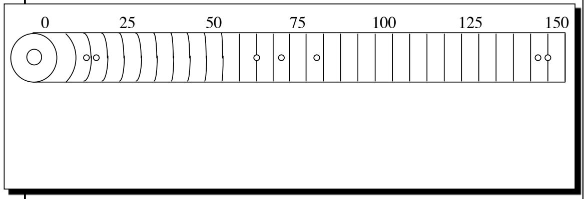
FCFS scheduling results in the head moving 550 tracks

Can we do better?

Disk Head Scheduling
Minimizing head movement

- Greedy scheduling: *shortest seek time first*
  - Rearrange queue from: 83 72 14 147 16 150

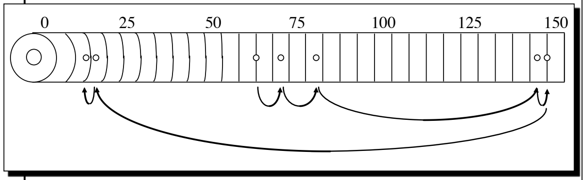
    To: 14 16 150 147 82 72



Disk Head Scheduling
Minimizing head movement

- Greedy scheduling: *shortest seek time first*
  - Rearrange queue from: 83 72 14 147 16 150

    To: 14 16 150 147 82 72

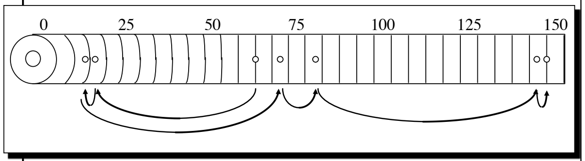


SSTF scheduling results in the head moving 221 tracks

Can we do better?

Disk Head Scheduling
SCAN scheduling

Rearrange queue from:	83	72	14	147	16	150
To:	150	147	83	72	14	16

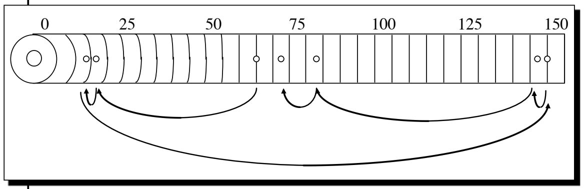


"*SCAN*" scheduling: Move the head in one direction until all requests have been serviced and then reverse. Also called elevator scheduling.

Moves the head 187 tracks

Disk Head Scheduling
Other variations

- *C-SCAN* scheduling ("*Circular*"-*SCAN*)
  - Move the head in one direction until an edge of the disk is reached and then reset to the opposite edge

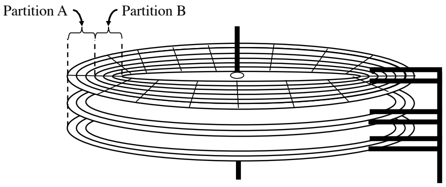


LOOK scheduling

Same as *C-SCAN* except the head is reset when no more requests exist between the current head position and the approaching edge of the disk

Disk Performance
Disk partitioning

- Disks are typically partitioned to minimize the largest possible seek time
  - A partition is a collection of cylinders
  - Each partition is a logically separate disk

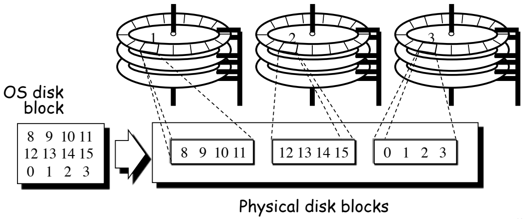


Disks – Technology Trends

- Disks are getting smaller in size
  - Smaller → spin faster; smaller distance for head to travel; and lighter weight
- Disks are getting denser
  - More bits/square inch → small disks with large capacities
- Disks are getting cheaper
  - 2x/year since 1991
- Disks are getting faster
  - Seek time, rotation latency: 5-10%/year (2-3x per decade)
  - Bandwidth: 20-30%/year (~10x per decade)
- Overall:
  - Disk capacities are improving much faster than performance

Management of Multiple Disks
Using multiple disks to increase disk throughput

- Disk striping (RAID-0)
  - Blocks broken into sub-blocks that are stored on separate disks
    - similar to memory interleaving
  - Provides for higher disk bandwidth through a larger effective block size



8	9	10	11
12	13	14	15
0	1	2	3

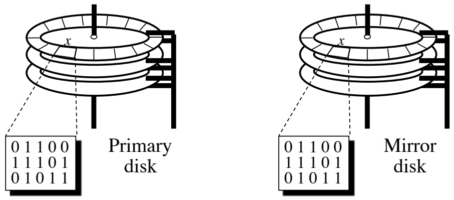
8 9 10 11

12 13 14 15

0 1 2 3

Management of Multiple Disks
Using multiple disks to improve reliability & availability

- To increase the reliability of the disk, redundancy must be introduced
  - Simple scheme: *disk mirroring (RAID-1)*
  - *Write to both disks, read from either.*



0 1 1 0 0
1 1 1 0 1
0 1 0 1 1

0 1 1 0 0
1 1 1 0 1
0 1 0 1 1

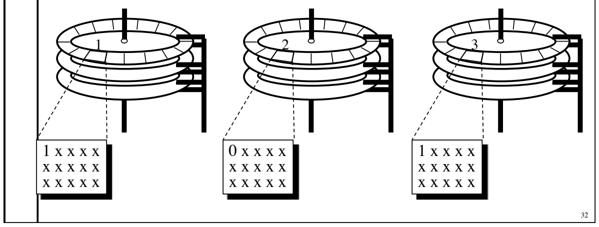
Who controls the RAID?

- Hardware
  - +Tend to be reliable (hardware implementers test)
  - +Offload parity computation from CPU
    - Hardware is a bit faster for rewrite intensive workloads
  - -Dependent on card for recovery (replacements?)
  - -Must buy card (for the PCI bus)
  - -Serial reconstruction of lost disk
- Software
  - -Software has bugs
  - -Ties up CPU to compute parity
  - +Other OS instances might be able to recover
  - +No additional cost
  - +Parallel reconstruction of lost disk

Management of Multiple Disks

Using multiple disks to increase disk throughput

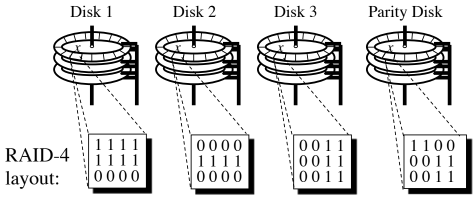
- RAID (*redundant array of inexpensive disks*) disks
  - Byte-wise striping of the disks (RAID-3) or block-wise striping of the disks (RAID-0/4/5)
  - Provides better performance and reliability
- Example: storing the byte-string 101 in a RAID-3 system



Improving Reliability and Availability

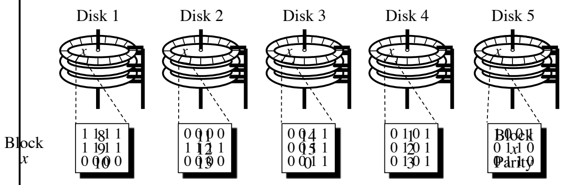
RAID-4

- Block interleaved parity striping
  - Allows one to recover from the crash of any one disk
  - Example: storing 8, 9, 10, 11, 12, 13, 14, 15, 0, 1, 2, 3



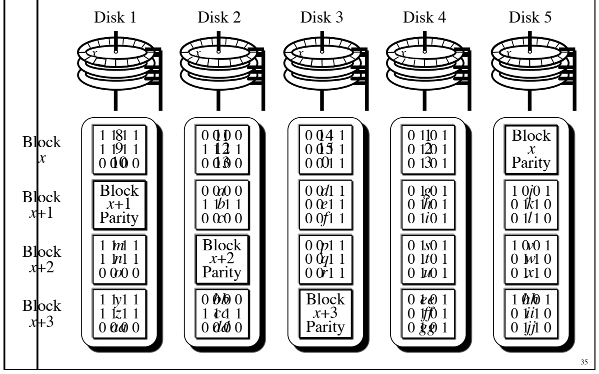
Improving Reliability and Availability

RAID-5 Block interleaved parity striping



Improving Reliability and Availability

RAID-5 Block interleaved parity striping


Disk 1
Disk 2
Disk 3
Disk 4
Disk 5
Block
x
Block
x+1
Block
x+2
Block
x+3
Block
x
Parity
Block
x+1
Parity
Block
x+2
Parity
Block
x+3
Parity